



# The Wireless FRINGE

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Advanced Architecture and Research

September 27<sup>th</sup>, 2013

# Agenda

The Fringe of the Internet

The Fringe backbone

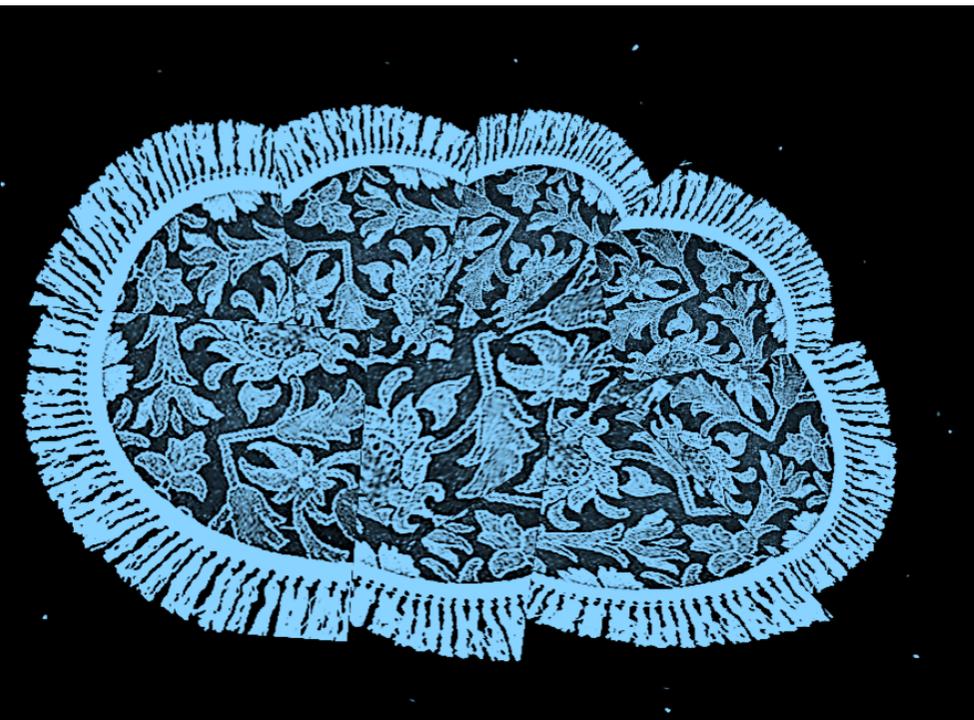
The Deterministic Fringe

6TSCH

Polymorphic Flows



# The routing infrastructure, today



## The Internet

Fully engineered

Hierarchical, Aggregations, ASs, Wire links

Fully distributed States

Shows limits (BGP tables, addr. depletion)

⇒ Reached adult size, mature to aging

⇒ Conceptually unchanged by IPv6

## IPv4 Intranets

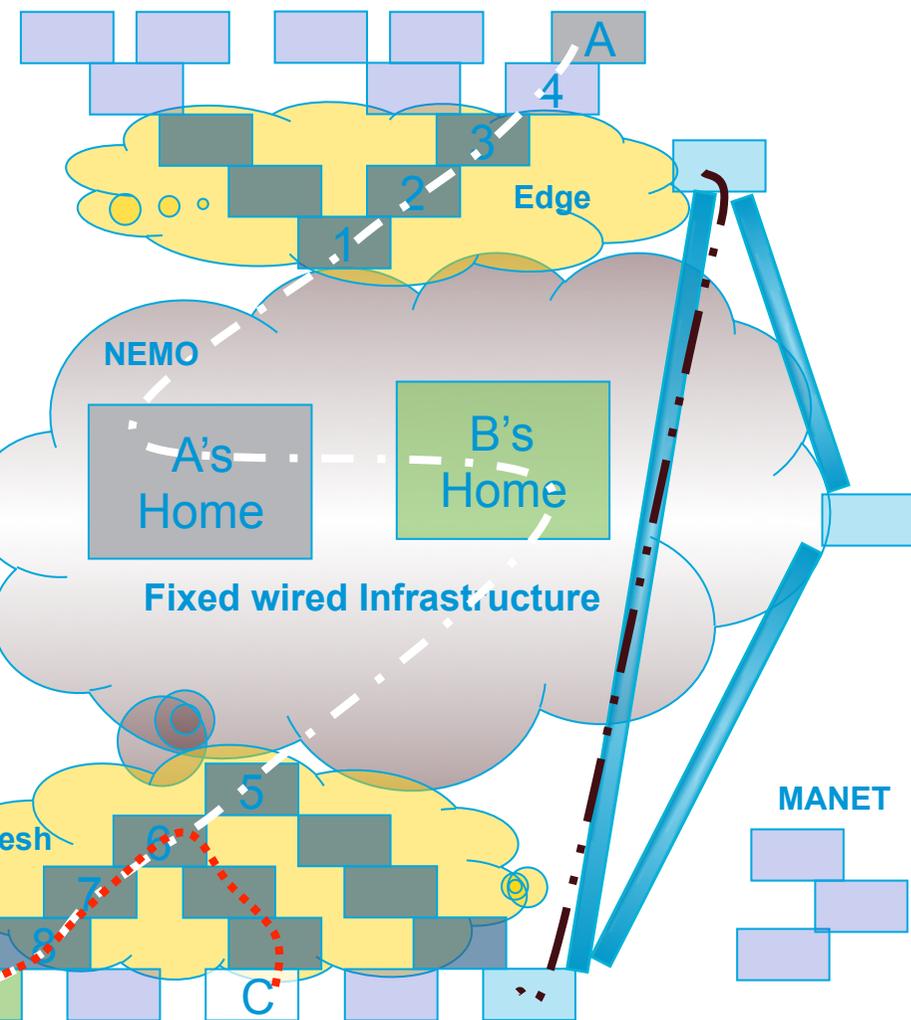
Same structure as the Internet

Yet decoupled from the Internet

NAT, Socks, Proxies

⇒ First model for Internet extension

# The emerging Fringe of the Internet



## L2 mesh Under

Multi-hop Public Access Points,  
Proprietary mission specific products  
Getting Deterministic

## L3 Route Over

Migration to IETF Protocols (RPL)  
Internet of Things (IOT, M2M)  
Different IPv6 (6LoWPAN, SDN)

## Mobile Overlays

Global reachability (NEMO, LISP, DMM)  
DataCenter Network virtualization (NVO3, LISP)  
Route Projection

**The Fringe DOES NOT LEAK  
into the  
Routing Infrastructure**

# A sense of history

Router only knows "self" with: ID, certificates  
Peers are discovered  
Links are discovered  
Routes are discovered  
=> Infinity of self-centric networks

## IPv6 autonomic

scale

## IPv4 routing

Router CLI with: ID, keys.  
All links to L2 peers  
Routes are discovered  
=> Single 'GRID'

## NA Subarea

NCP generation with:  
All Transmission Groups to L2 peers  
All Physical Units type 4 nodes,  
All Virtual Routes

# Scaling to Pervasive IoE/T

000\*<sup>scale</sup> => No leak in the Internet

=> Opaque Fringe operations

=> Federating backbone

Reachability

=> Radio

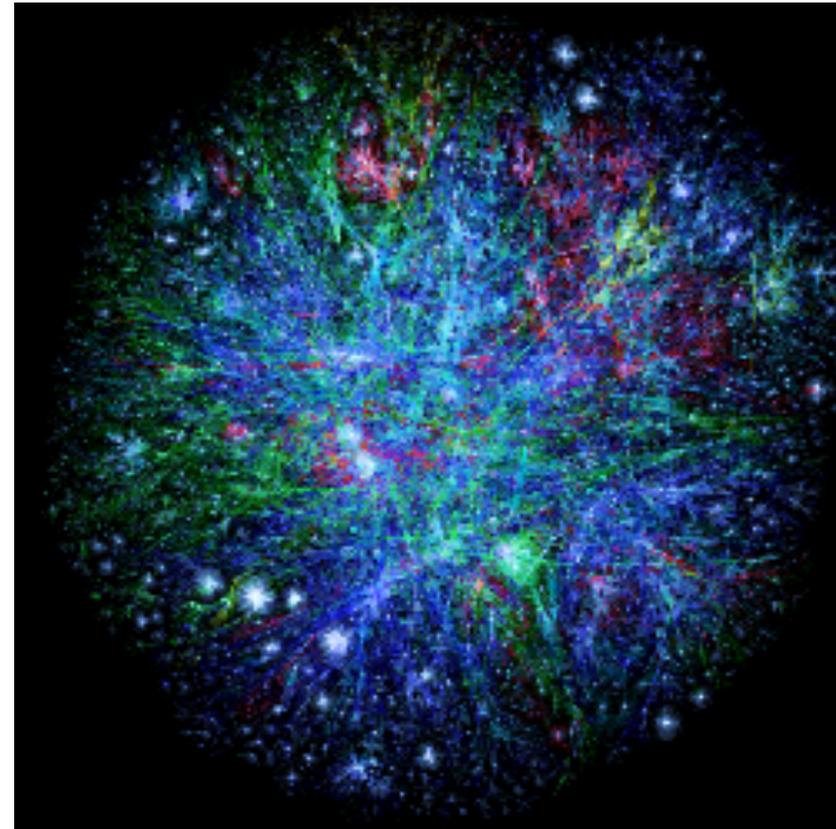
Addressing

=> IPv6

Density

=> spatial reuse

=> Routing



# Agenda

✓ The Fringe of the Internet

The Fringe backbone

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Polymorphic Flows



# Backbone Devices

Wired or Wireless

Switches and Controllers

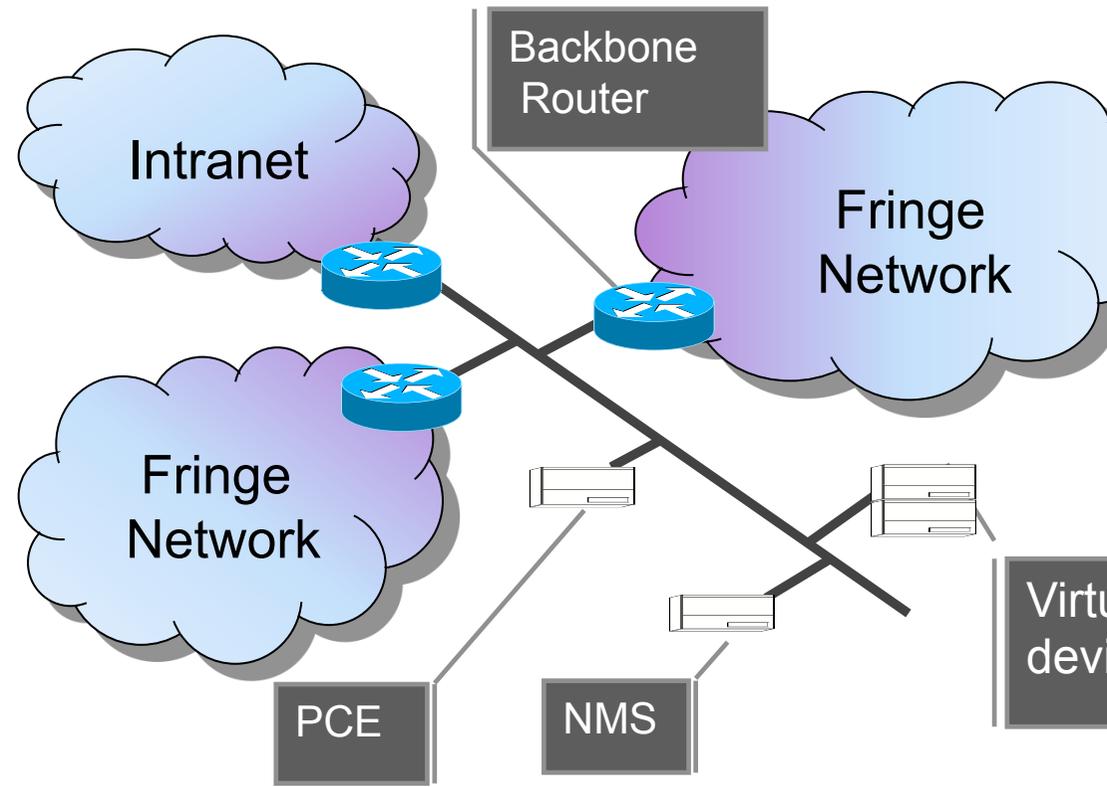
Wireless ND

Global Mobility integration

Virtualized Functions

PCE for deterministic route computation

Legacy Appliances e.g. industrial PLC



# The IPv6 Neighbor Discovery (R)evolution

A new Efficient ND, aka WiND for Wireless ND

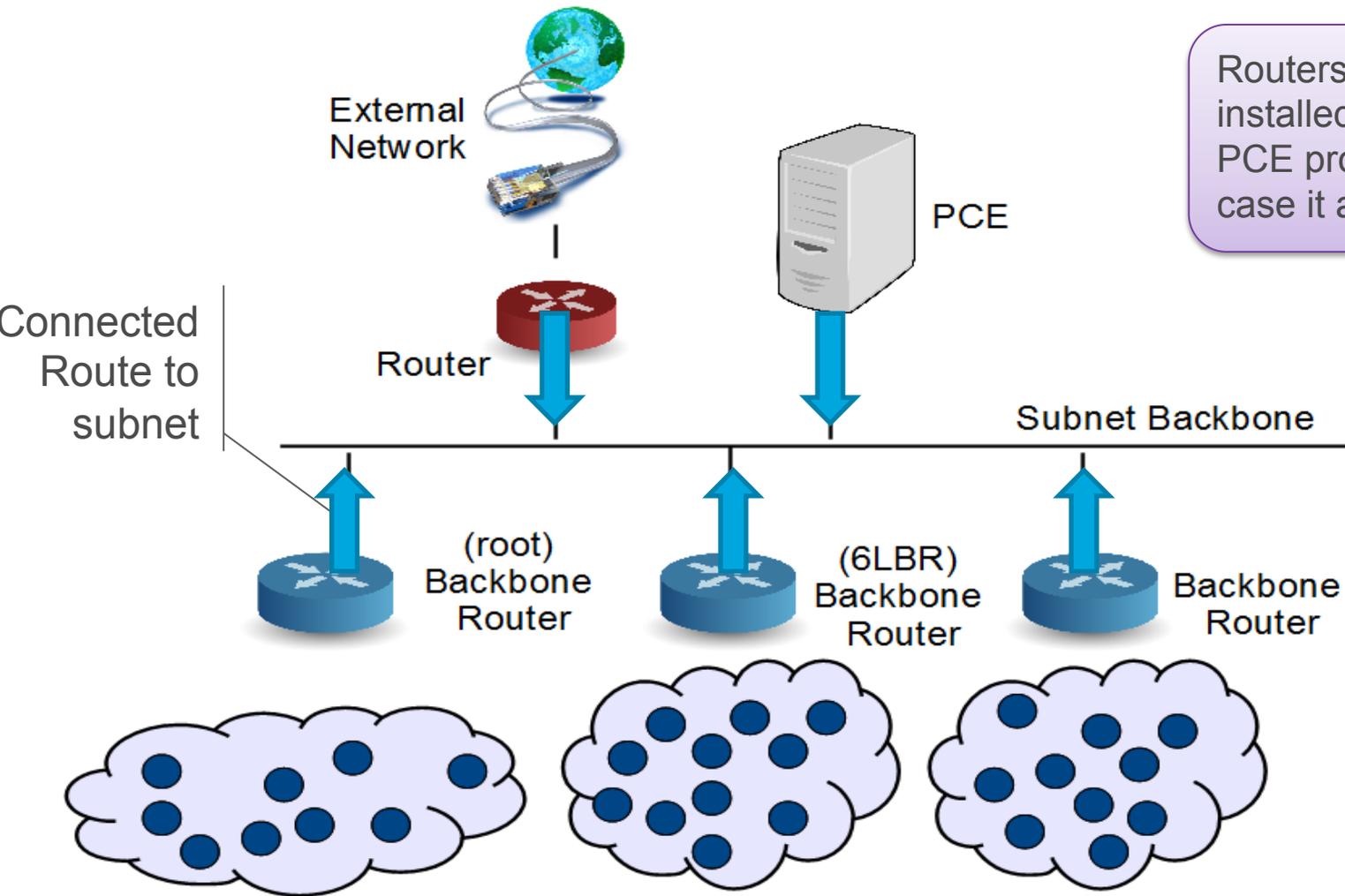
Registration and Duplicate Address Detection (DAD)  
extending 6LoWPAN ND over the backbone

Resolution

Distributed through proxy-ND

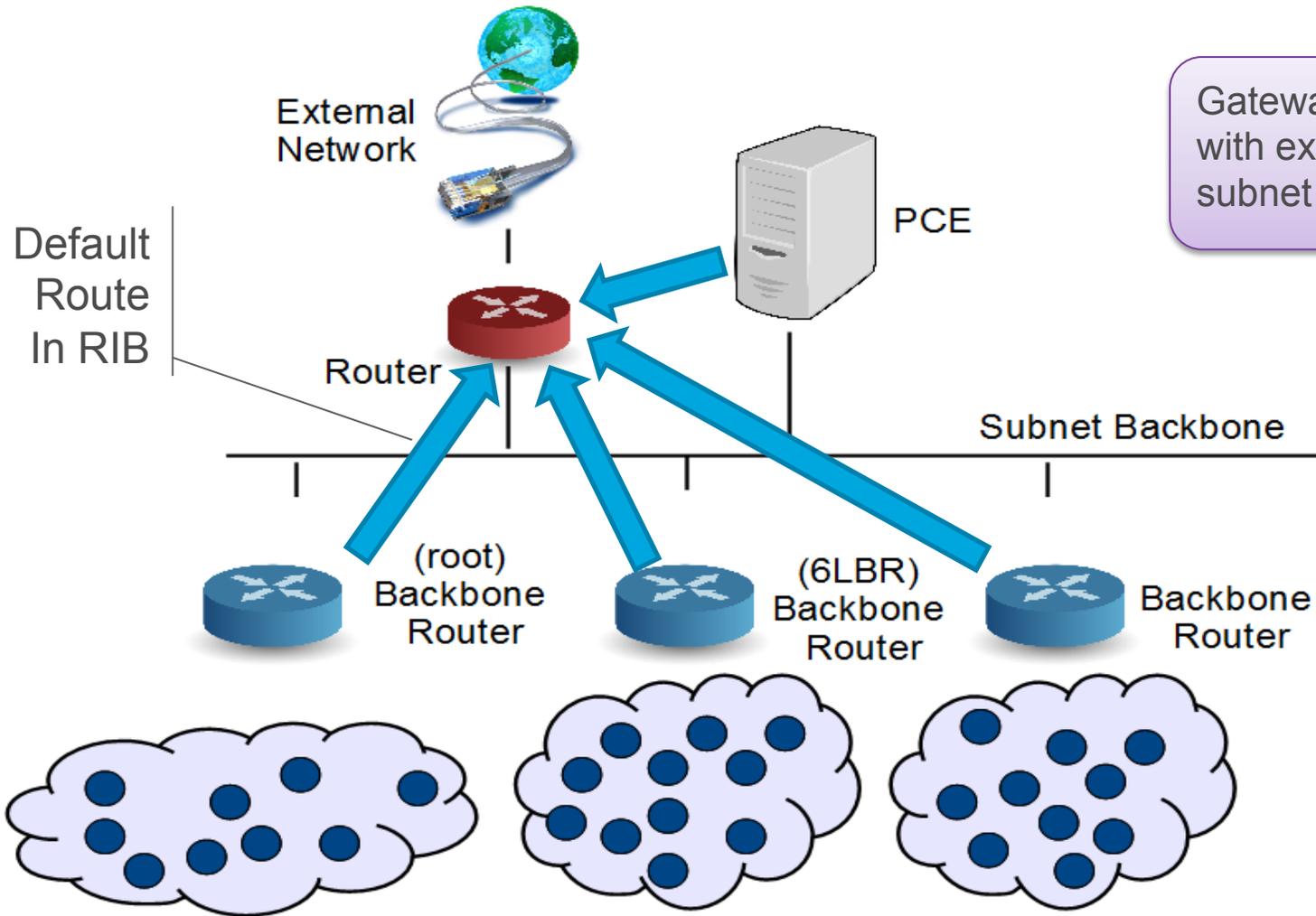
Routing in not-onlink mode

# Initial time



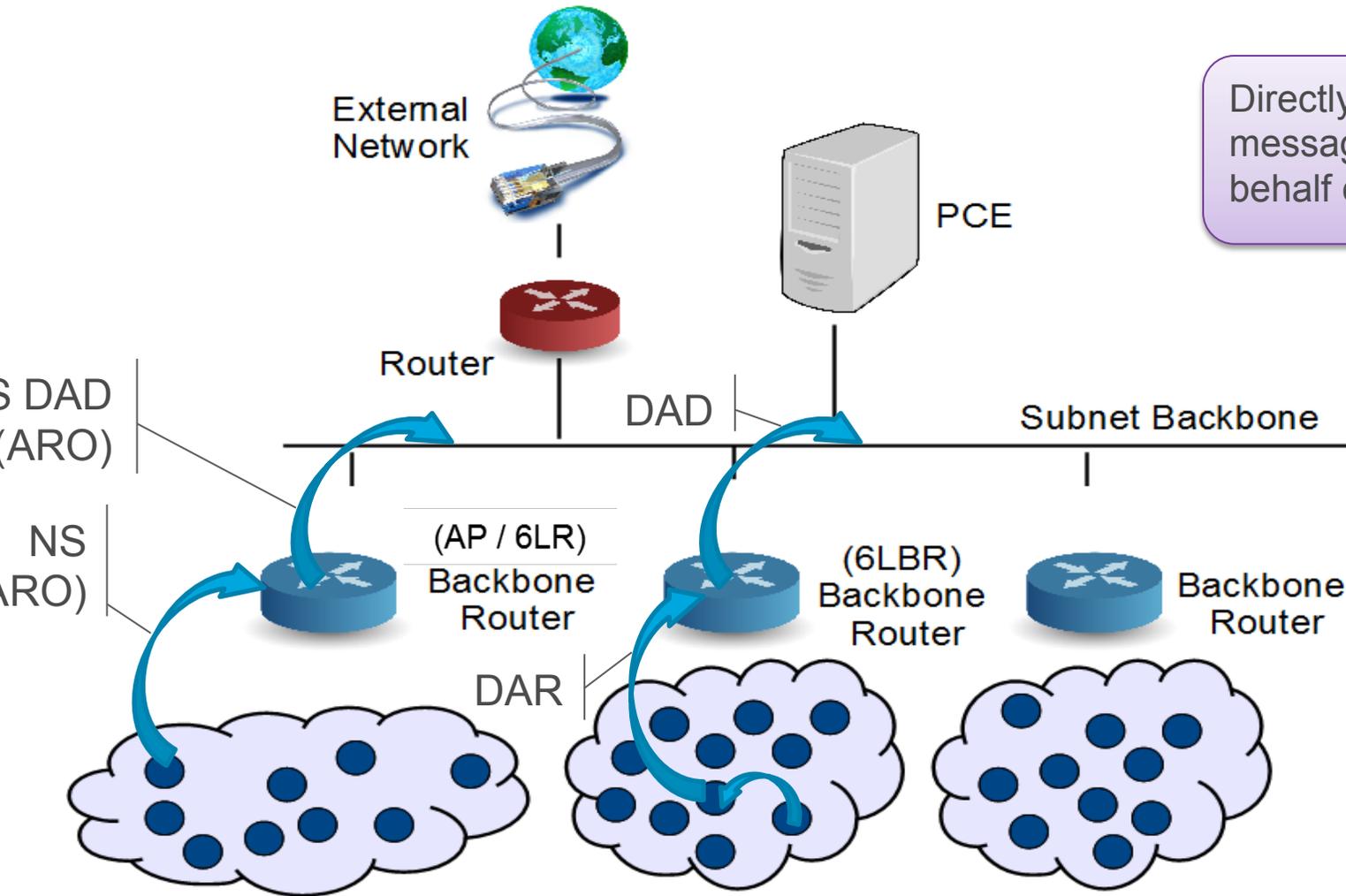
Routers within subnet have a connected route installed over the subnet backbone.  
PCE probably has a static address in which case it also has a connected route

# First advertisements from GW (RA, IGP, RPL)



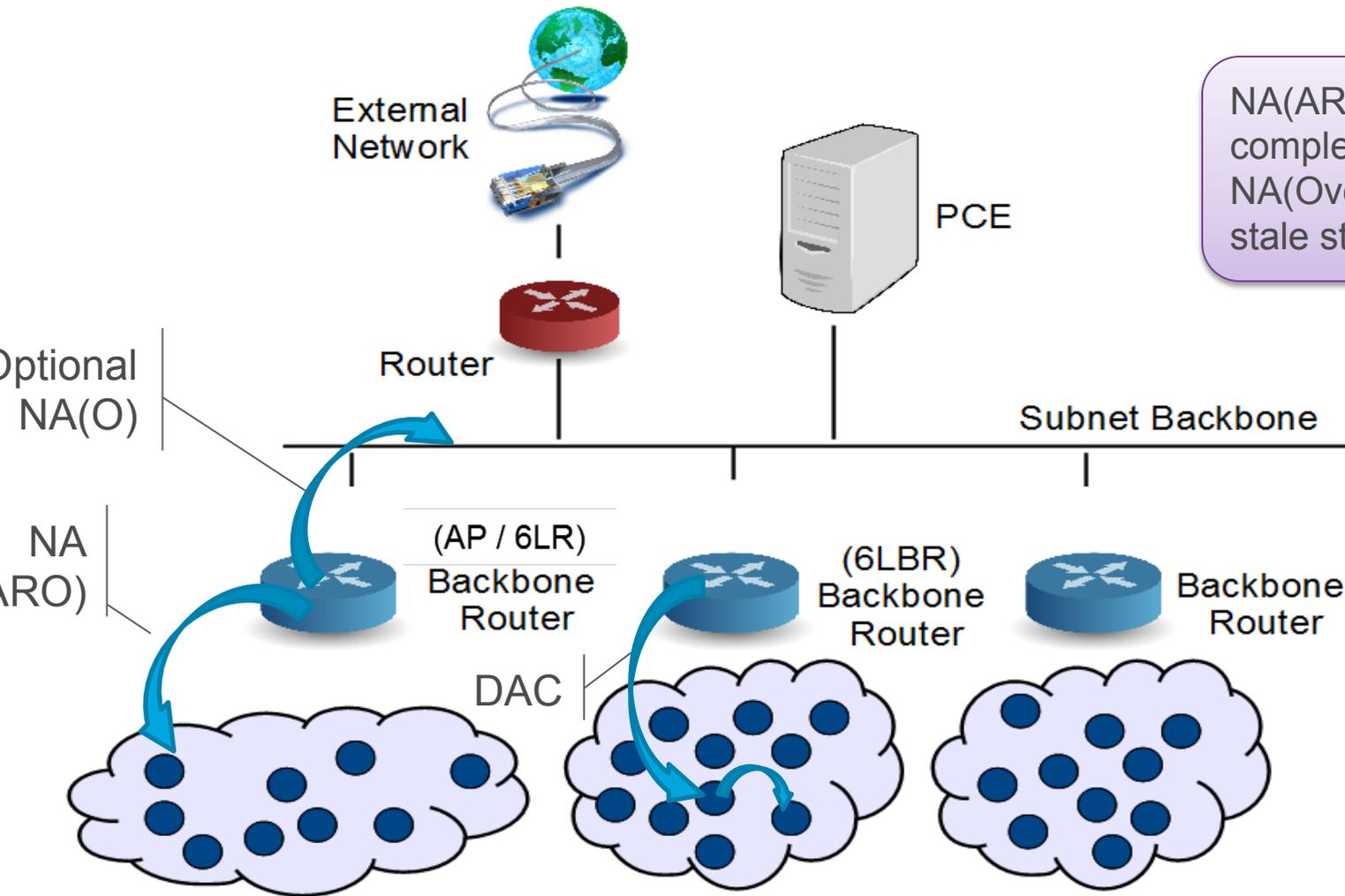
Gateway to the outside participate to some IGP with external network and attracts all extra-subnet traffic via protocols over the backbone

# IPv6 ND Registration to 6LR and 6LBR



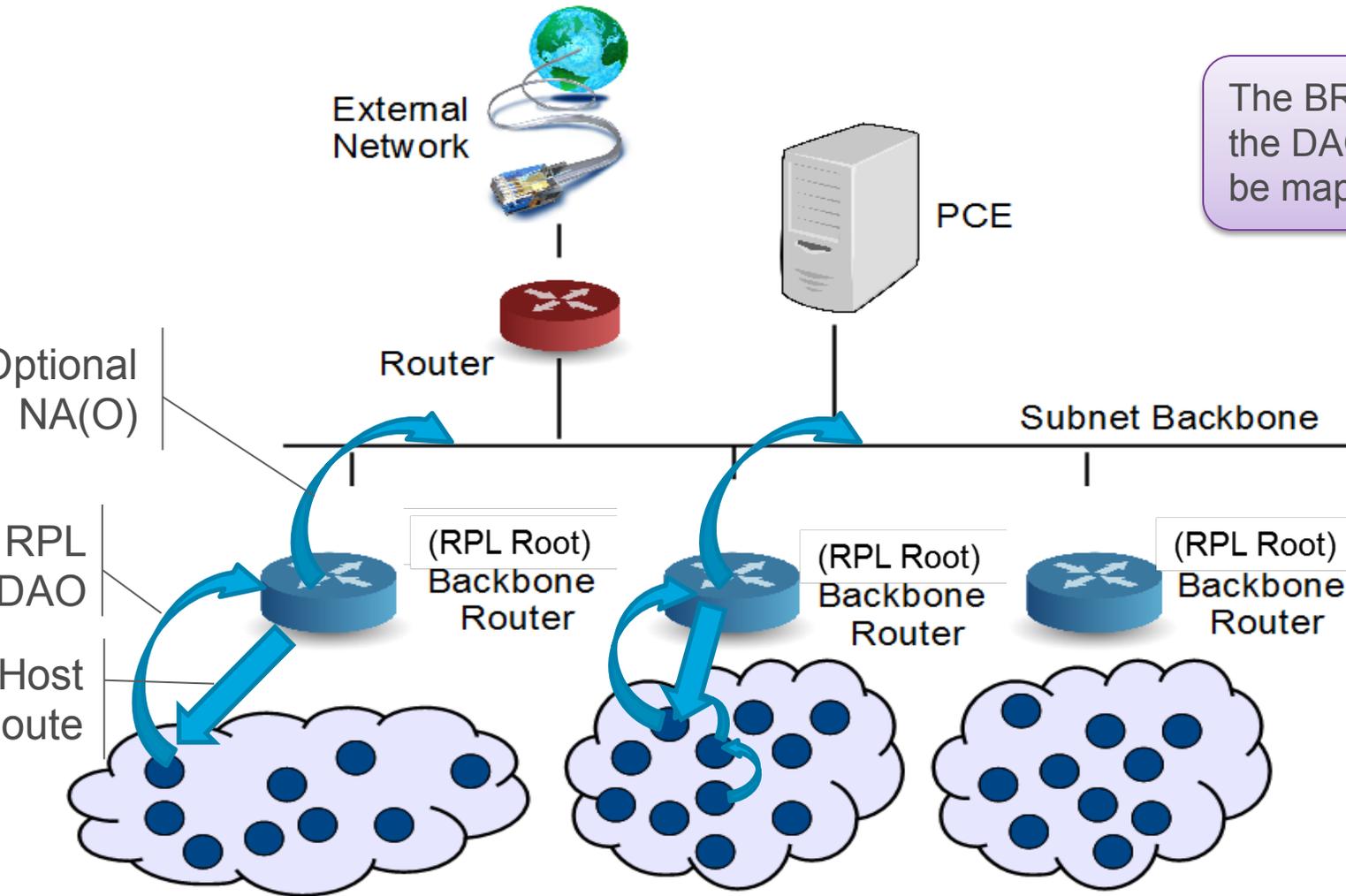
Directly upon NS(ARO) or indirectly upon DA message, the backbone router performs DAD behalf of the wireless device.

# IPv6 ND Registration and Proxy for NS ARO



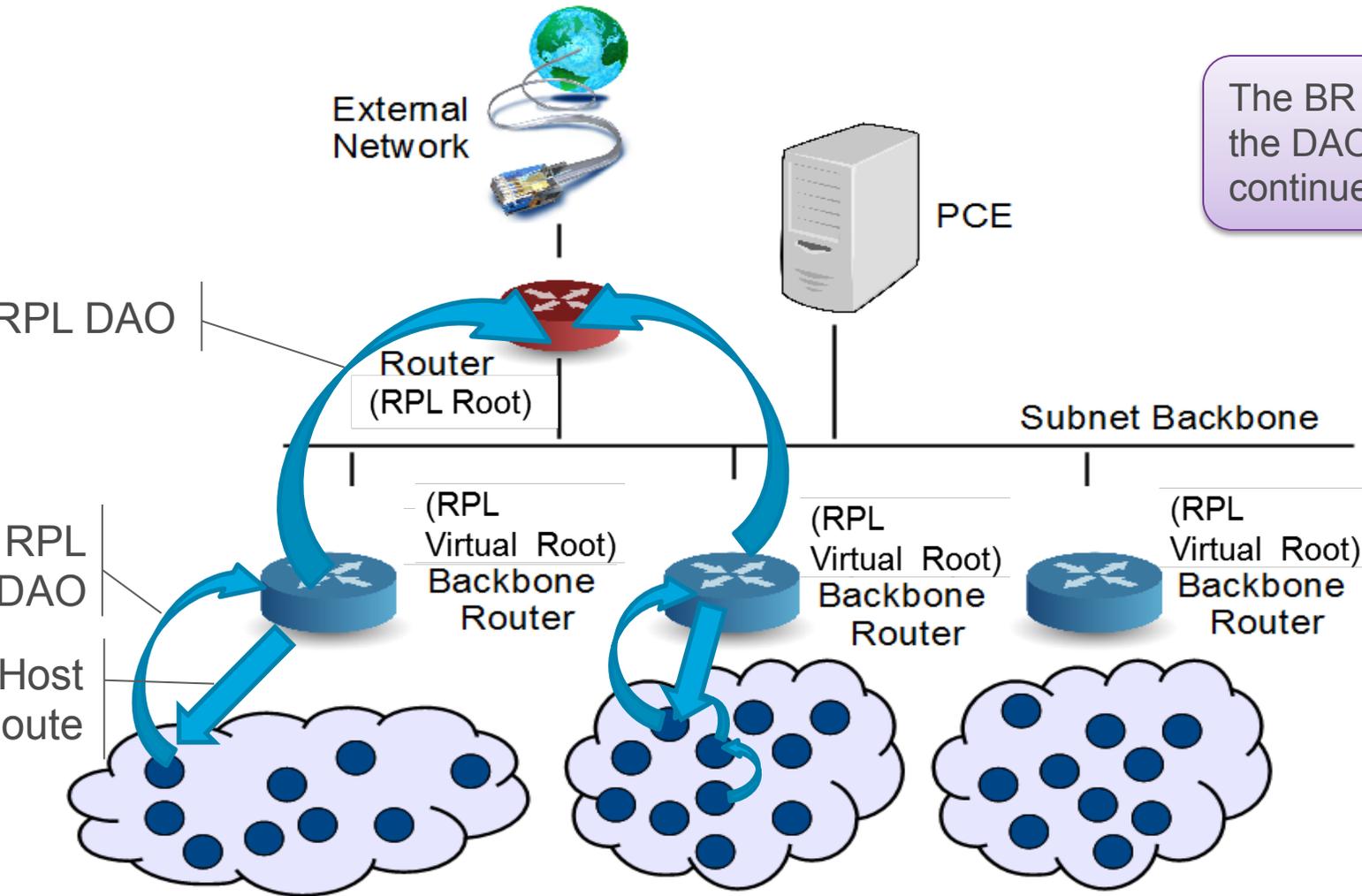
NA(ARO) or DAC message carry successful completion if DAD times out.  
NA(Override) is optional to clean up ND cache stale states, e.g. if node moved.

# IPv6 ND Proxy for RPL



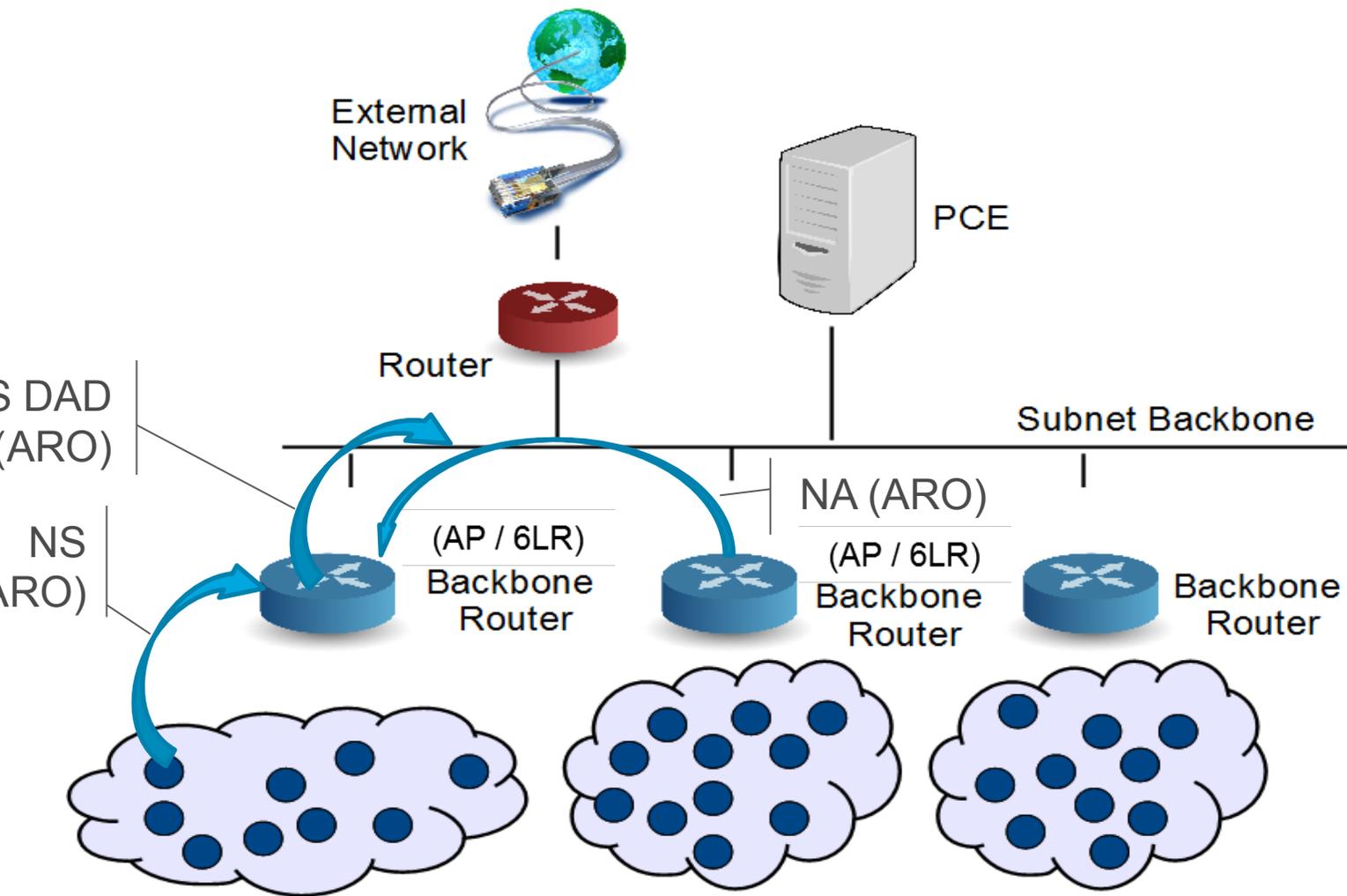
The BR maintains a route to the WSN node for the DAO Lifetime over instance VRF. VFR may be mapped onto a VLAN on the backbone.

# RPL over the backbone



The BR maintains a route to the WSN node for the DAO Lifetime over instance VRF that is continued with RPL over backbone.

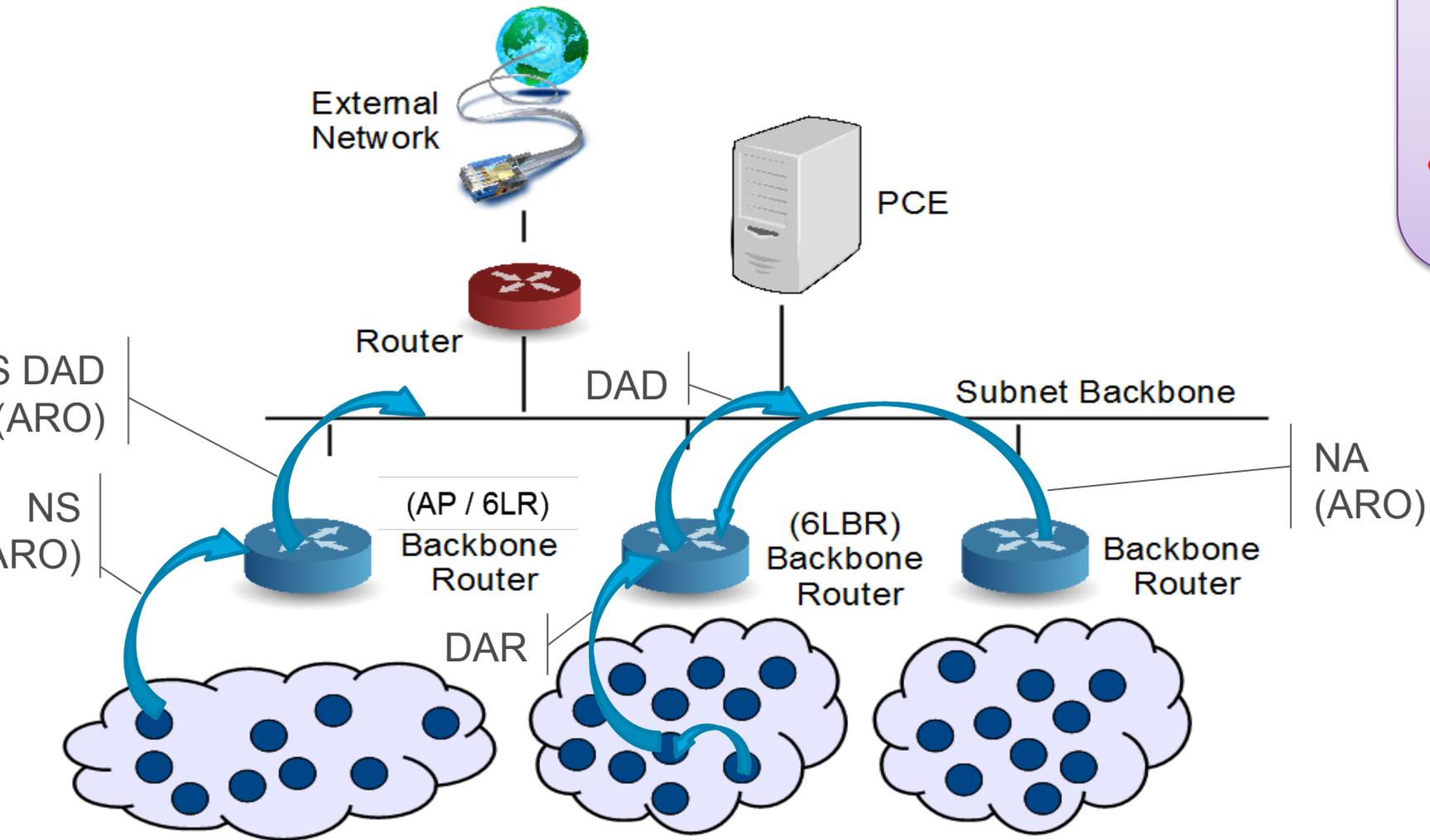
# Duplication



DAD option has:  
Unique ID  
TID (SeqNum)

Defend with NA if:  
**Different OUID**  
Newer TID

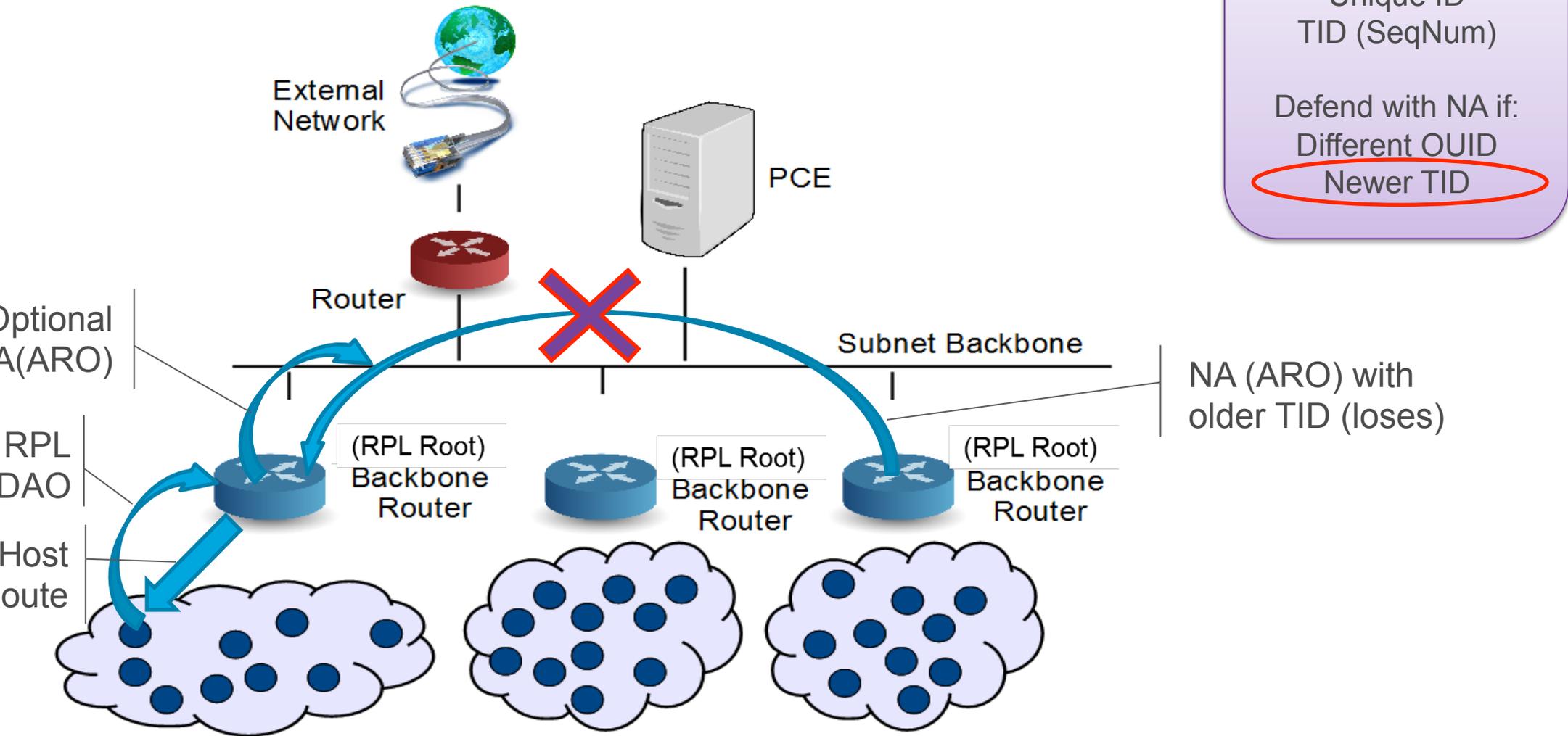
# Mobility



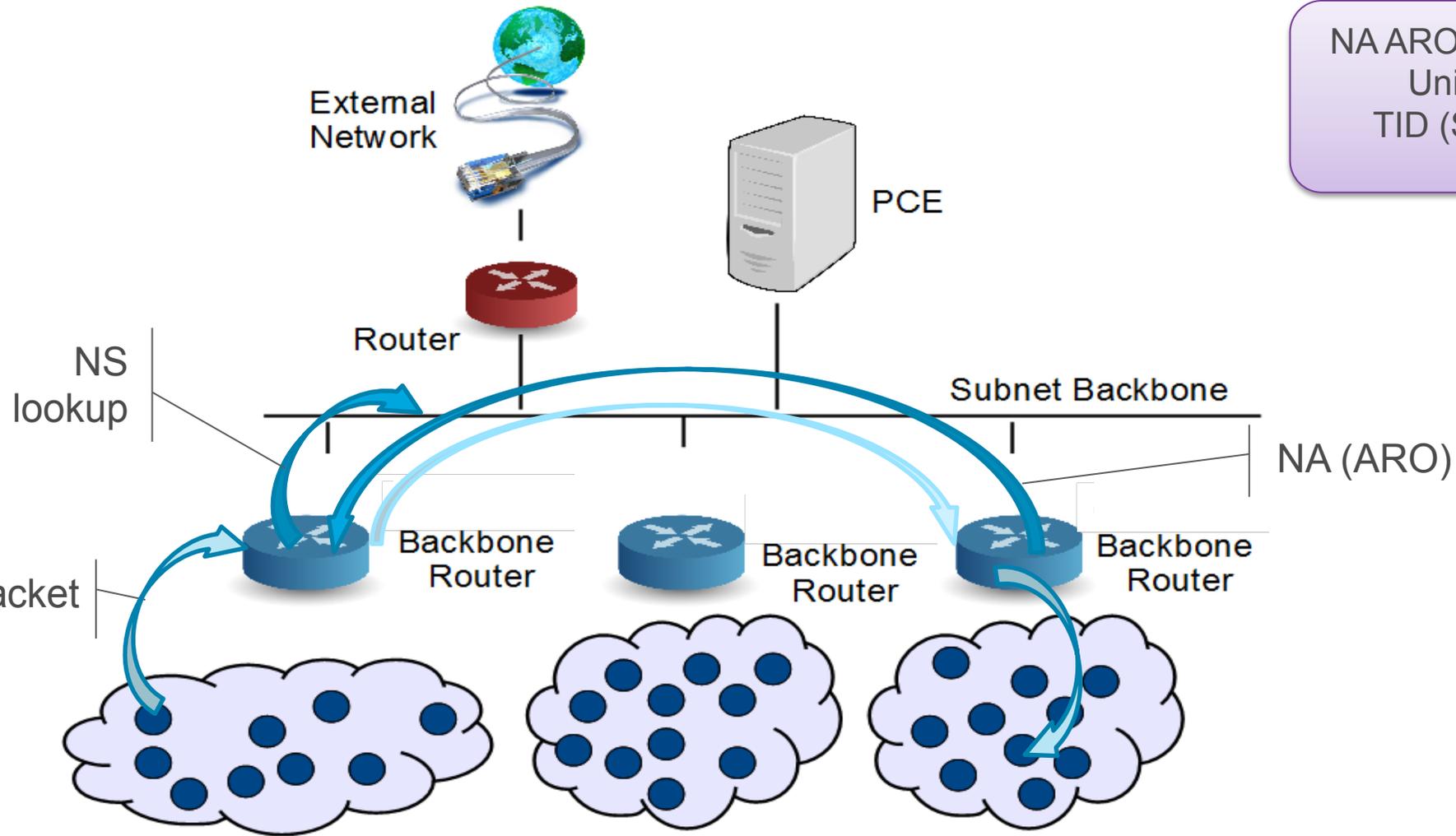
DAD option has:  
Unique ID  
TID (SeqNum)

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Newer TID

# Mobility

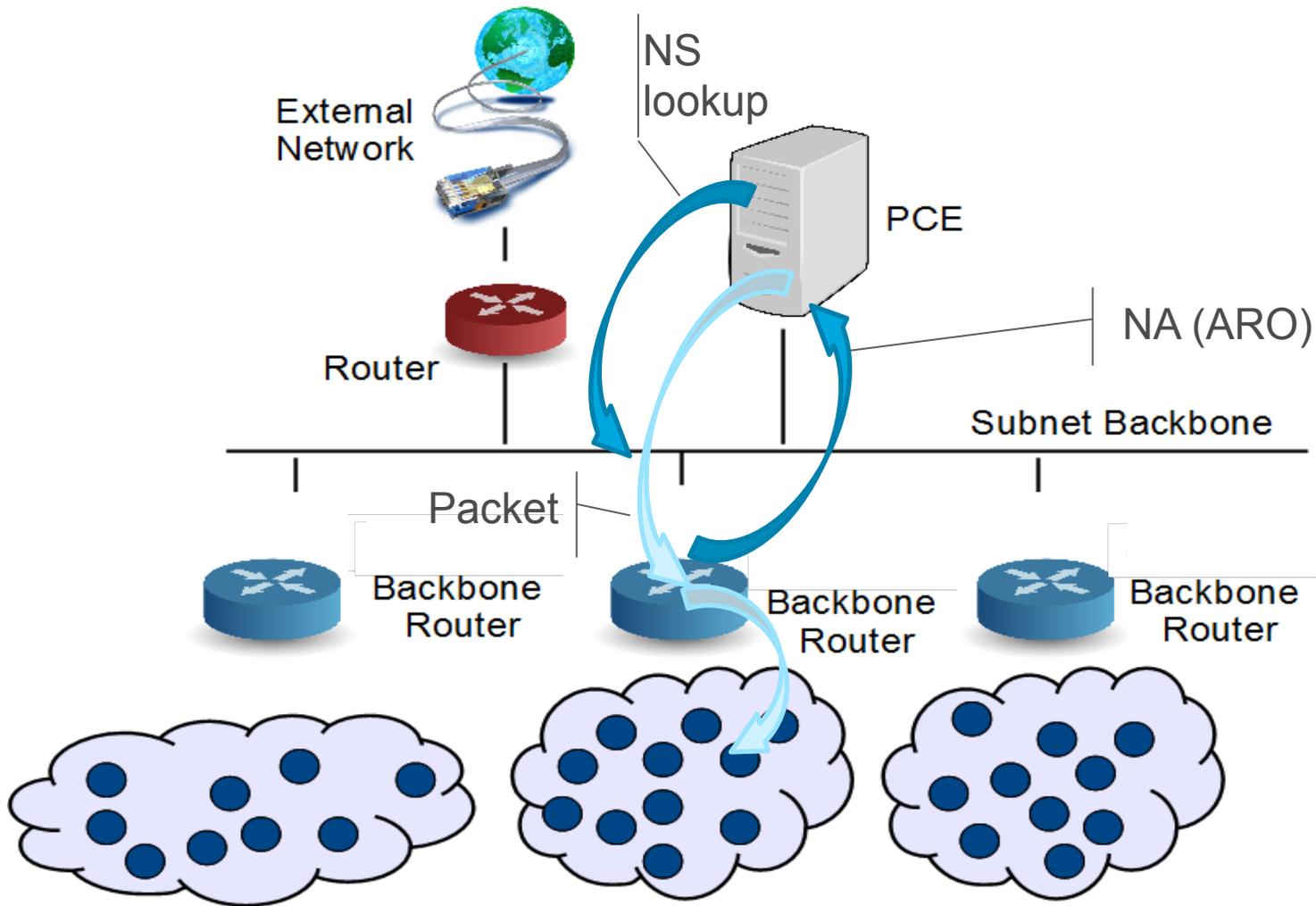


# Resolution



NA ARO option has:  
Unique ID  
TID (SeqNum)

# Resolution (2)



Mixed mode ND  
BBR proxying over  
the backbone

# Agenda

✓ The Fringe of the Internet

✓ The Fringe backbone

The Deterministic Fringe

6TSCH

Polymorphic Flows



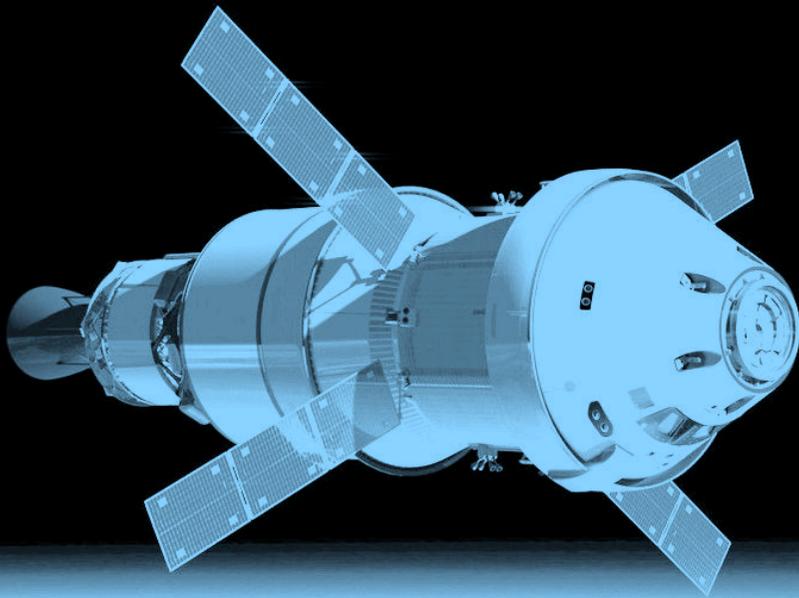
# Deterministic Networking

new level of guarantees

For deterministic traffic (known a priori)

Network synchronization and Timely transmission

Centralized routing and scheduling optimization



Cleveland, Columbus & Cincinnati R

### SPECIAL TIME SCHEDULE

FOR THE TRAIN CONVEYING THE  
REMAINS OF ABRAHAM LINCOLN, LATE PRESIDENT OF THE U. S., AND  
FROM WASHINGTON, D. C., TO SPRINGFIELD, ILL.

Cleveland to Columbus, Saturday, April 29th, 1865

Leave Cleveland.....	12.00	Midnight.
Berea.....	12.43	A. M.
Olmsted.....	12.51	"
Columbia.....	1.02	"
Grafton.....	1.23	"
La Grange.....	1.37	"
Wellington.....	2.00	"
Rochester.....	2.17	"
New London.....	2.36	"
Greenwich.....	2.59	"
Shiloh.....	3.10	"
Shelby.....	3.39	"
Crestline.....	4.07	"
Galion.....	4.23	"
Iberia.....	4.41	"
Gilend.....	5.05	"
Cardington.....	5.20	"
Ashley.....	5.43	"
Eden.....	5.55	"
Berlin.....	6.19	"
Lewis Centre.....	6.32	"
Orange.....	6.47	"
Worthington.....	6.56	"
Arrive Columbus.....	7.30	A. M.

This Train will have exclusive right to the Road against all other  
Public Locomotive will be run ten minutes in advance of the  
regular time.

E. S. FLINT, Superintendent

# Process Control example: a refinery

Sensors and actuator usually close: <100m

Elements of critical loops can be wire-interconnected

Low-quality (noisy, intermittent) field power often available

Control room typically 500m to 2km distant

Wired: power + signal carried together on one twisted pair

Designed to meet intrinsic safety regulations: ~40mW/pair

Typically 25-pair to 100-pair wire bundles in buried conduit

Wireless is most attractive for this long, costly link

Primary requirement is  $\geq 5$ yr battery life for field devices

Battery replacement is often very costly or impractical

Environmental power harvesting strongly desired

e.g. ExxonMobil Baytown refinery near Houston (Texas)

100 hectares (40 sq. miles) with 12 control rooms (CRs): 10 refining CRs localized to a few hectares each,

2 CRs that span the plant and interact with the other 10; large chemical plant adjacent to refinery



# What's specific

Industrial networks evolved concurrent with IP

Process Control and Factory Automation goals

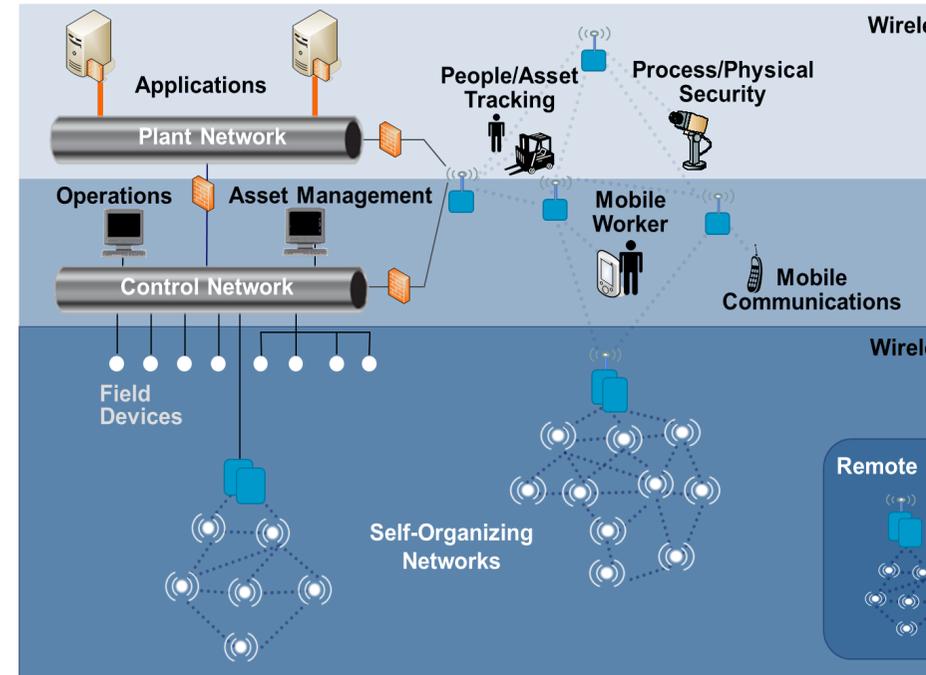
differ from those of IT => schism

Need low latency and high cyclic determinism

Need higher reliability, faster repair, lesser maintenance

Largest plants potentially reap the most profit from small productivity improvements

thus more willing to innovate (in stages) and to fund proven new technology



# Agenda

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- ✓ The Fringe backbone
- ✓ The Deterministic challenge

6TSCH

Polymorphic Flows



# 6TSCH: standardizing IPv6 over TSCH MAC

Align existing standards

(RPL, 6LoWPAN, PANA, RSVP, PCEP, MPLS) over 802.15.4e TSCH

Support Mix of centralized and distributed deterministic routing

Design 6TUS sublayer for L3 interactions

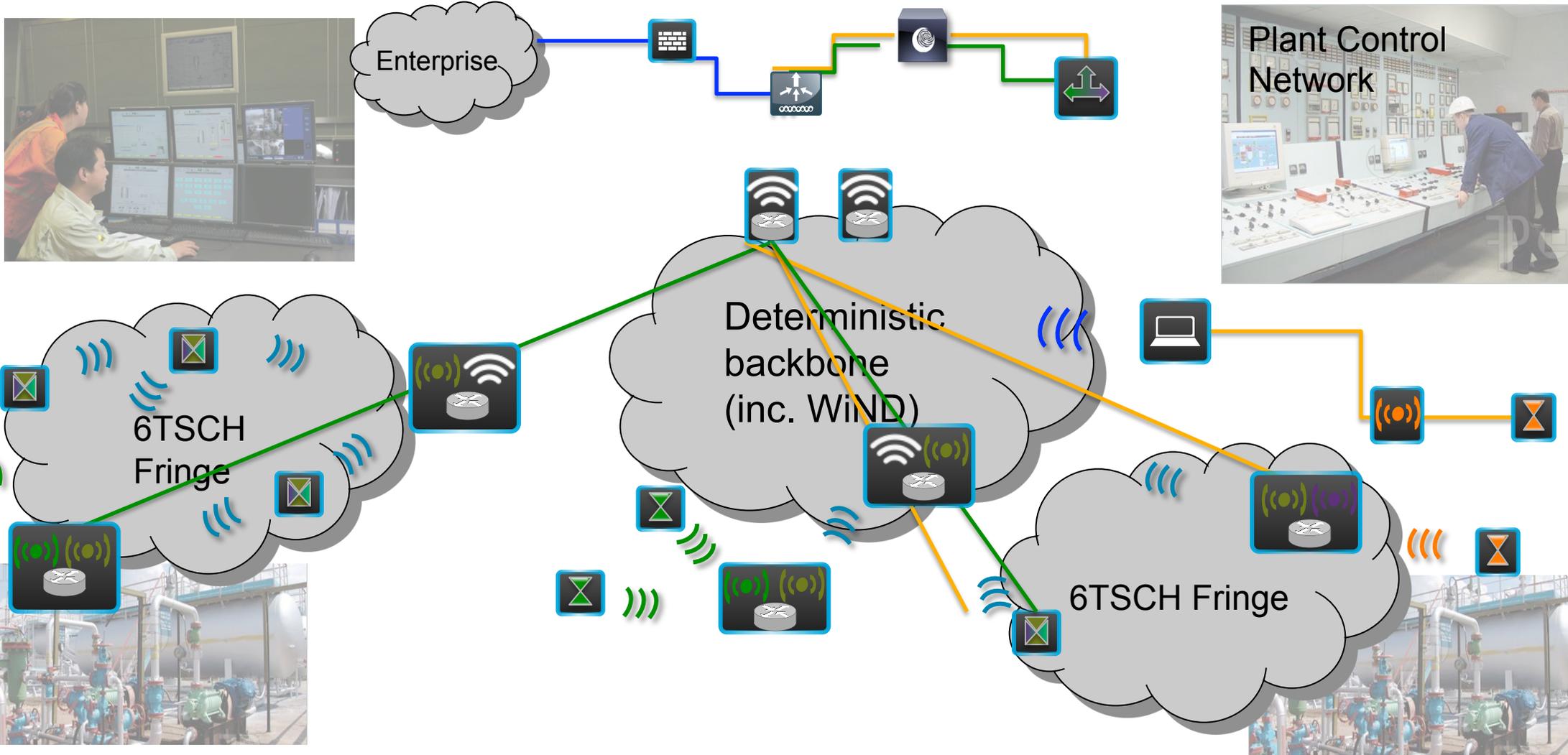
Open source implementations (openWSN...)

Build Ecosystem (DUST, NIVIS, IoT6, uni.lux, uni.bari, Berkeley...)

6 active drafts (Archi, 6TUS, ...)

Preparing for BoF in Berlin this summer

# Future Industrial Architecture with unified wireless



# TSCH: Architecture

Path  
Computation  
Element

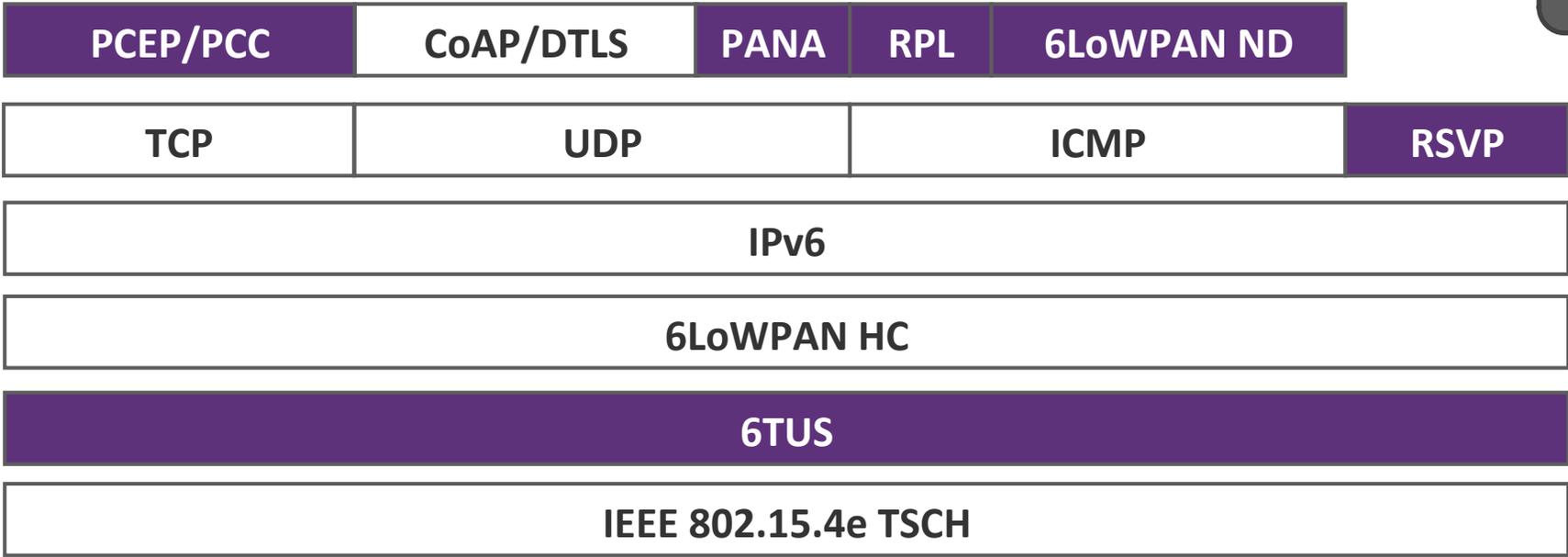
Datagram  
Transport  
Layer  
Security

Protocol for carrying  
Authentication for  
Network Access

6LoWPAN Neighbor Discovery

Resource Reservation  
Setup Protocol

Compares to  
ISA100.11a  
Data Link Layer



# Centralized vs. Distributed routing

Centralized

Distributed

God's view optimization

Autonomic & Mobile

Multipath redundancy

Highly available (DARPA)

Deterministic (optimized)

Deterministic

Virtualization

Scalability

# Routing With RPL

## Low Power Lossy Nets

Dynamic Topologies

Peer selection

Constrained Objects

Fuzzy Links

Routing, local Mobility

Global Mobility

## Addressed in RPL ?



Distance Vector + stretch



Peer only with parents



DV + Non-storing mode



Lazy Update & datapath valid

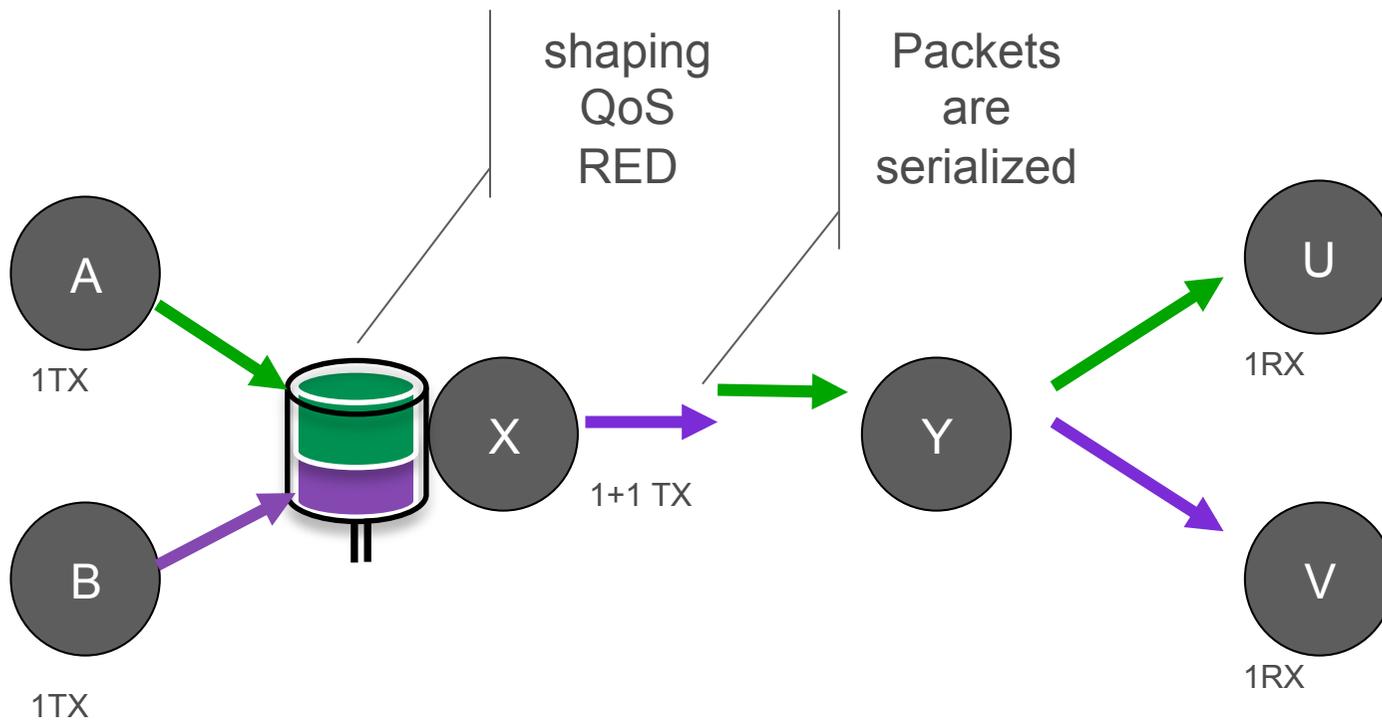


Constrained instances, TID

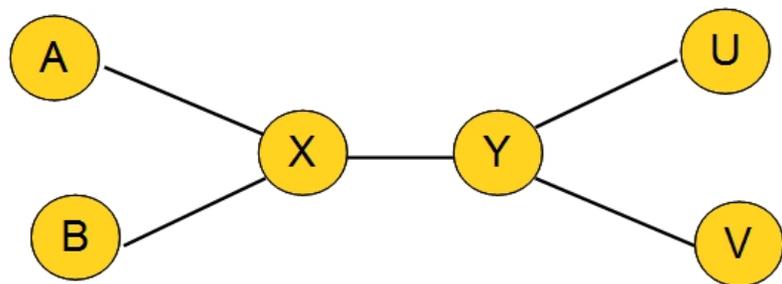


Req coupling with LISP/NEMO

# Normal L3 operation



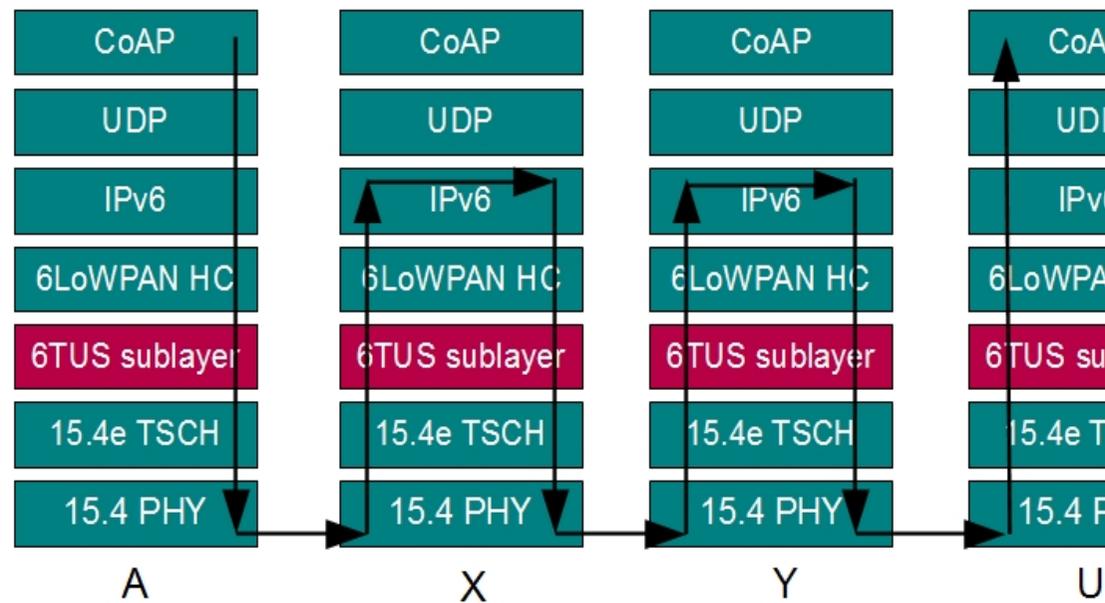
# Best effort routing



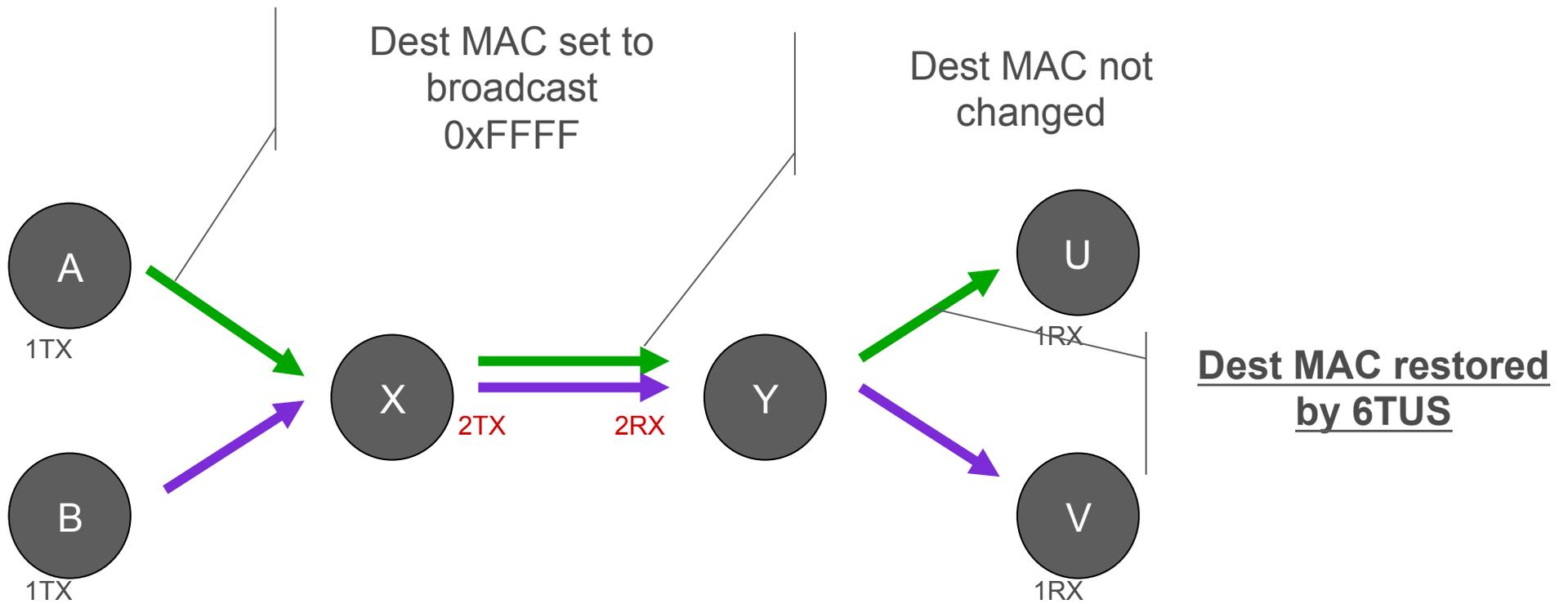
Bundle

			X → Y			Y → V
A → X					X → Y	Y → U
	B → X					

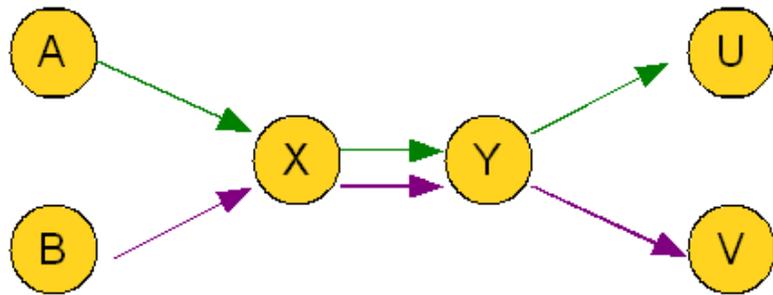
slotOffset



# Normal Track operation



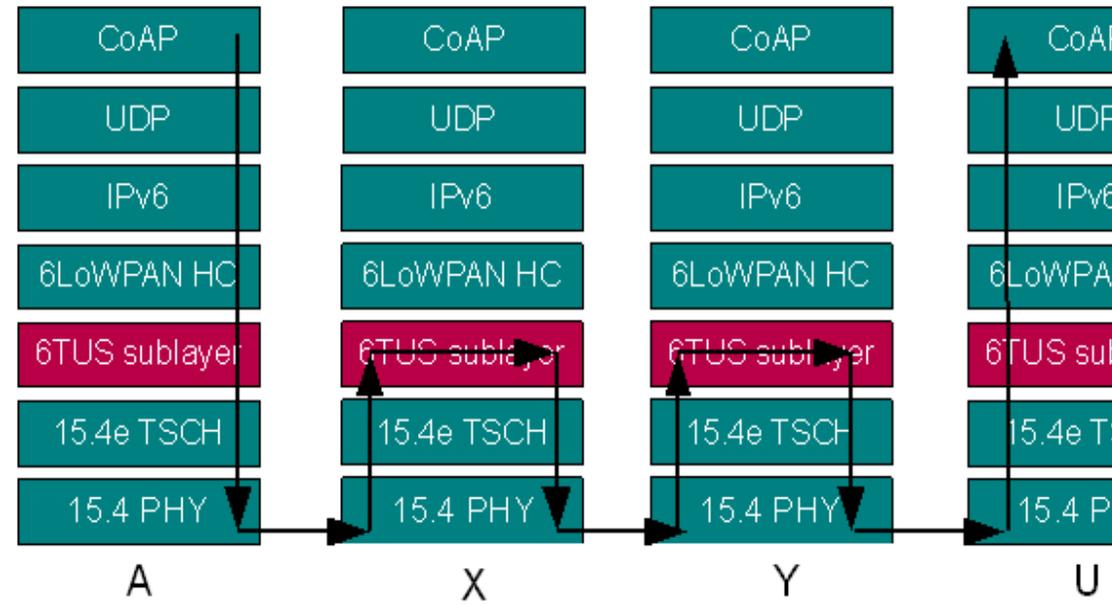
# Track Switching (G-MPLS)



Track

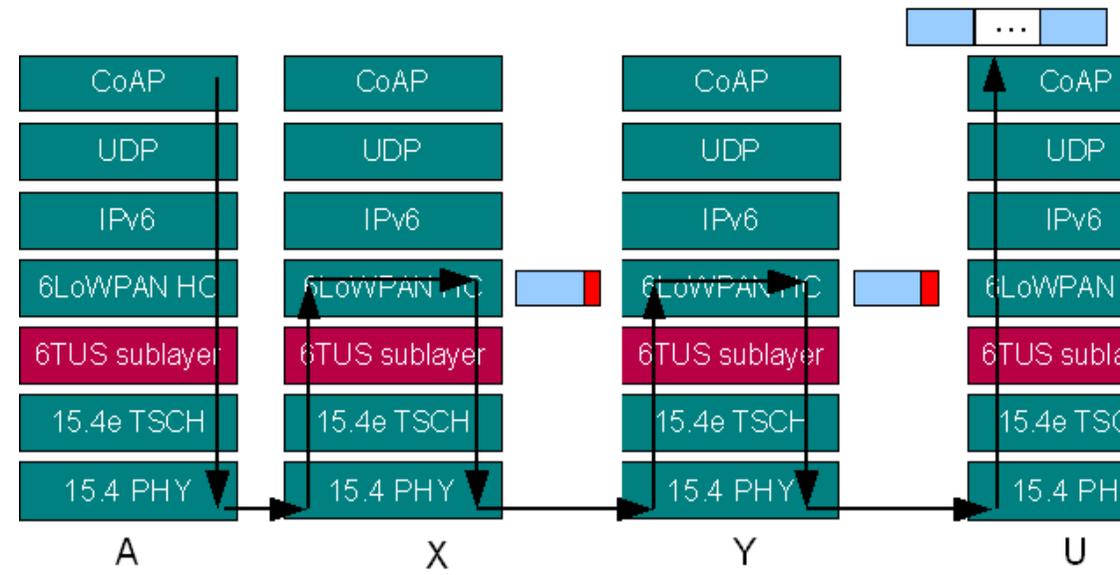
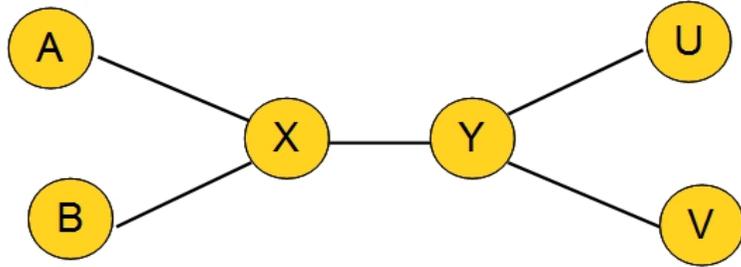
			X → Y			Y → V	
		X → Y		X → Y			
A → X					X → Y		Y → U
	B → X						

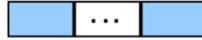
slotOffset





# 6LoWPAN Fragment forwarding



 packet, composed by N fragments  
 i-th fragment of the packet, with  $i > 1$   
 state installed in a single fragment

Bundle

			X → Y		Y → V	
A → X				X → Y		Y → U
	B → X					

slotOffset

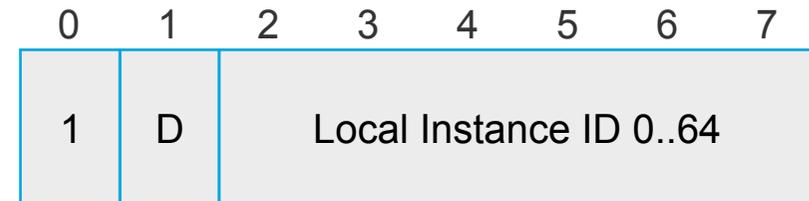
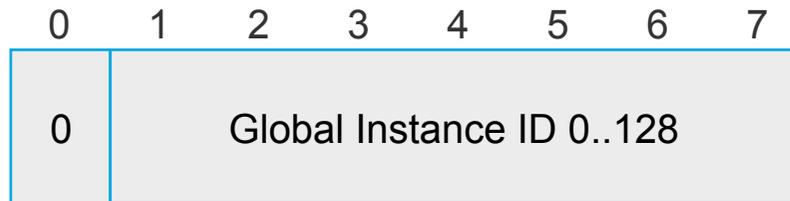
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- ✓ The Fringe of the Internet
- ✓ The Fringe backbone
- ✓ The Deterministic challenge
- 6TSCH
- Polymorphic Flows



# RFC 6550: RPL Instance ID

The RPL instance ID allows different routing optimizations, constraints and policies.



The RPL instance ID is encoded in 1 octet. The first bit indicates whether Global or Local.

“A local RPLInstanceID is autoconfigured by the node that owns the DODAGID and it MUST be unique for that DODAGID. The DODAGID used to configure the local RPLInstanceID MUST be a reachable IPv6 address of the node, and it MUST be used as an endpoint of all communications within that Local instance.”

Inside a packet: “If the 'D' flag is set to 1, then the destination address of the IPv6 packet MUST be the DODAGID. If the 'D' flag is cleared, then the source address of the IPv6 packet MUST be the DODAGID.”

6TSCH extends RPL's language of DODAGID to route (reservation) endpoint.

# Global Instances

28 global instances per network

Indexed by tuple (IPv6, InstanceID)

Running as Ships-in-the-night

1 instance = 1 VRF = 1 « L3 vlan »

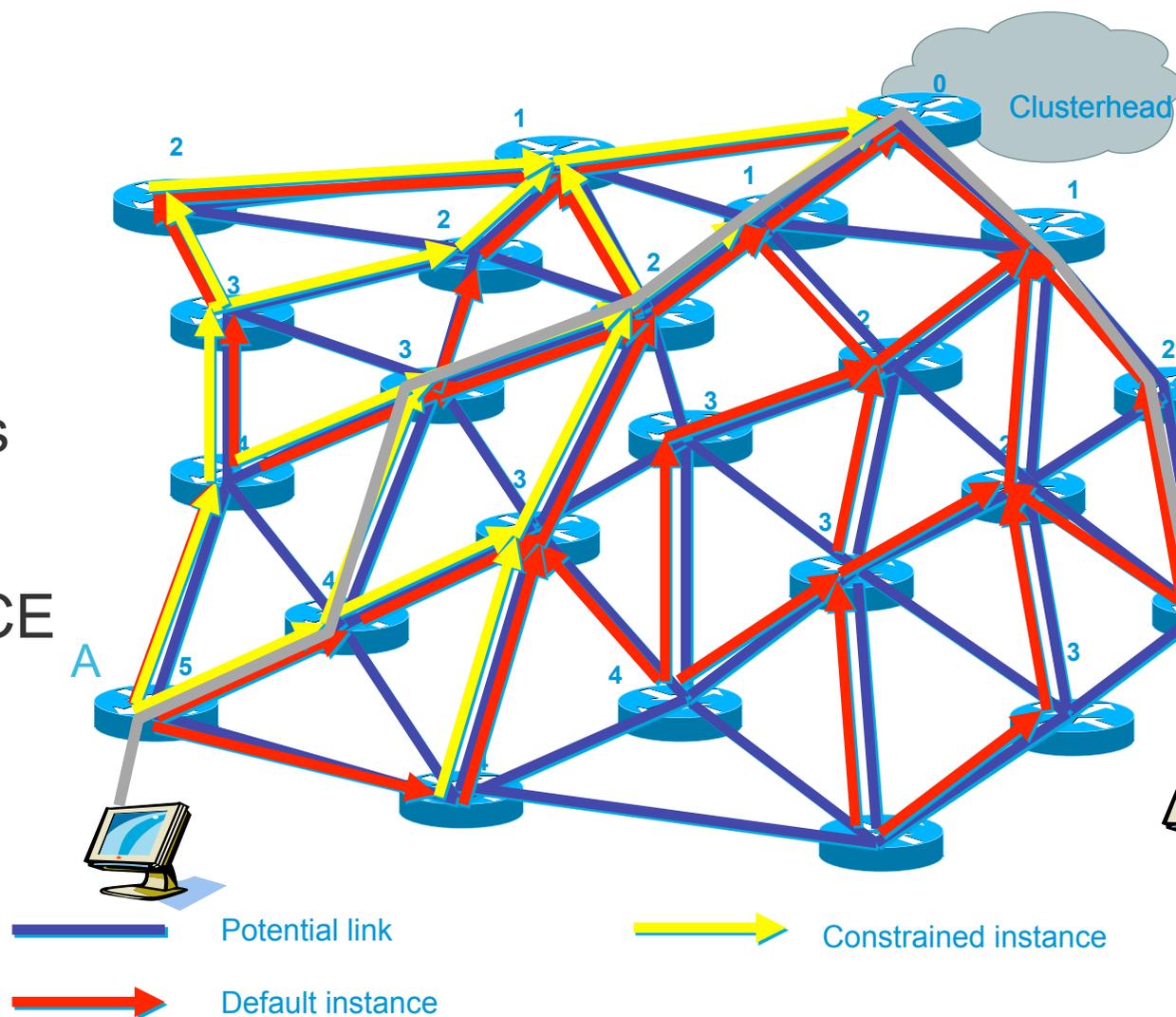
Serving different Objective Functions

Using different metrics

Can be shared between RPL and PCE

RPL along a DODAG

PCE adding orthogonal shortcuts



# Local instances

## 64 local Instances

per IPv6 source address

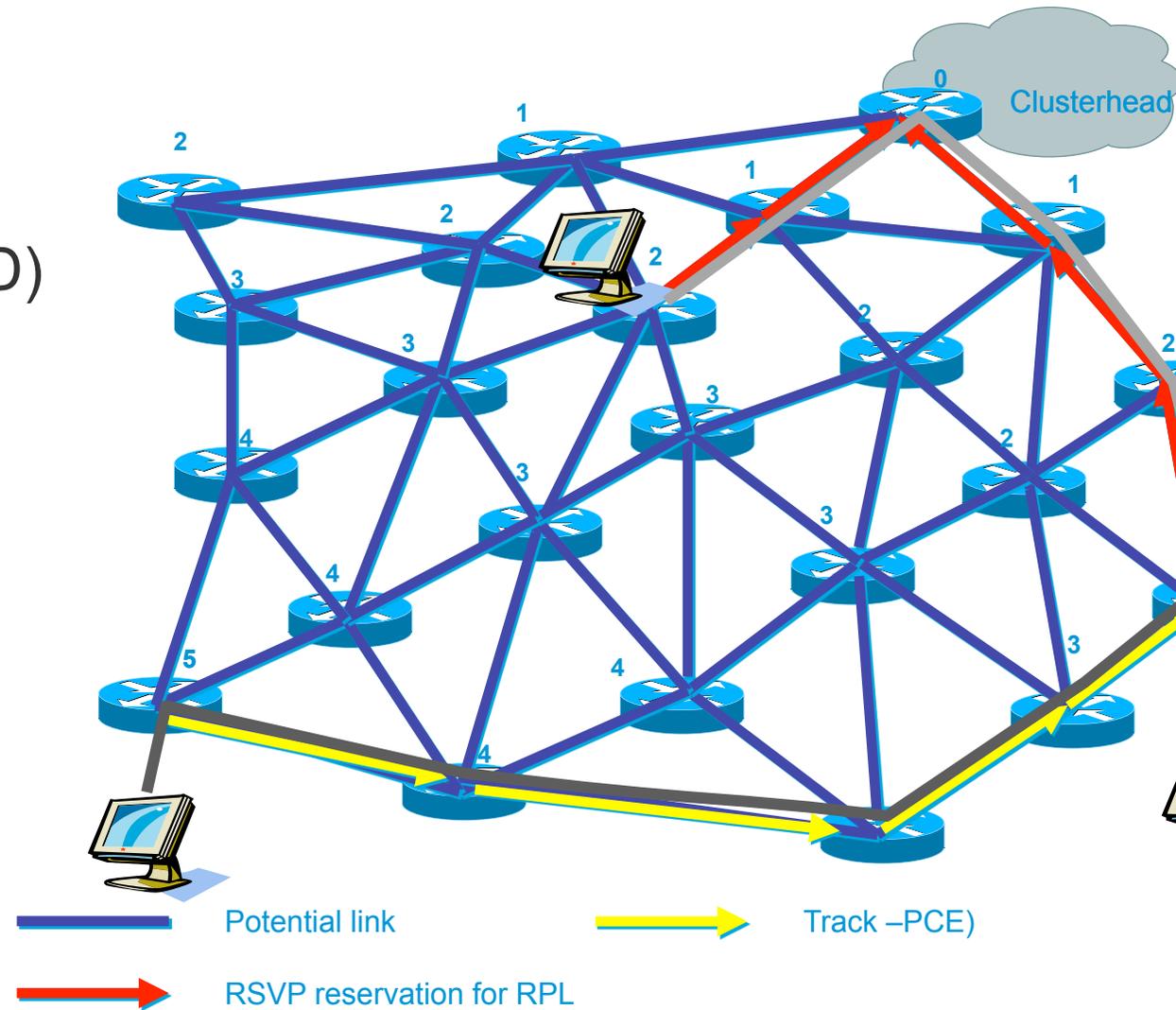
Indexed by tuple (IPv6, InstanceID)

## Used by RPL or PCE

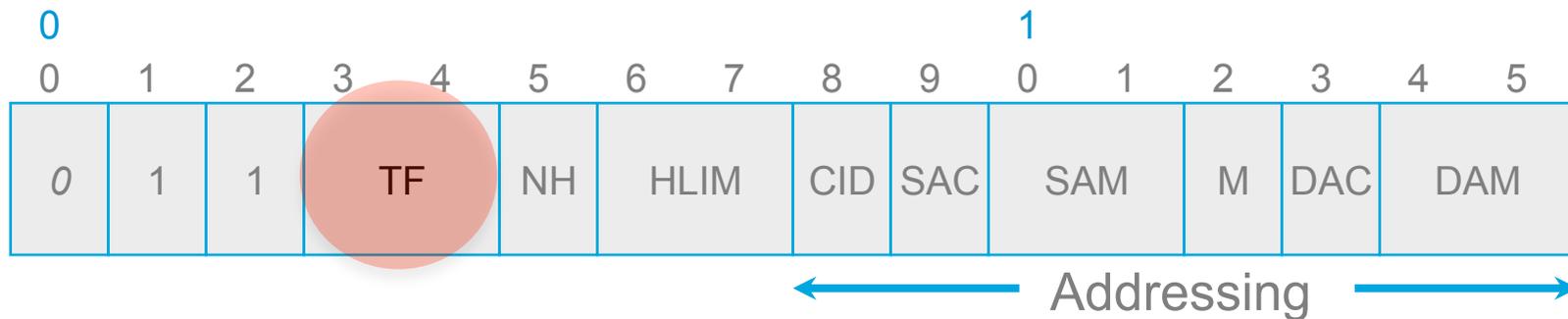
RPL: for P2P applications

RPL: to index RSVP path

PCE: Serves as Track ID,  
included in PCEP request  
from the source device

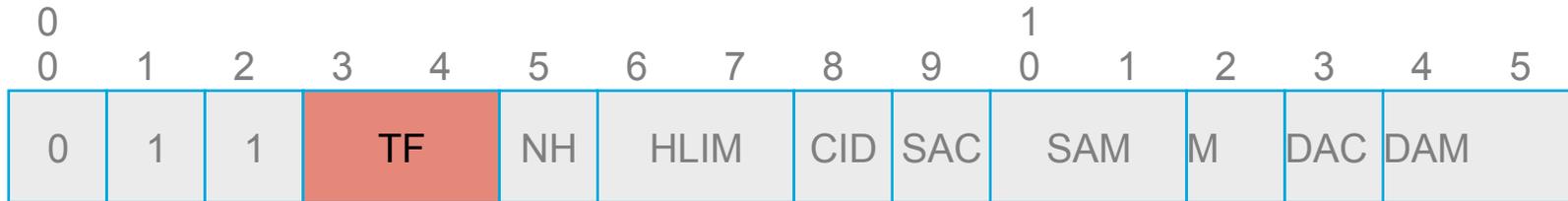


# RFC 6282: 6LoWPAN IPv6 Header Compression

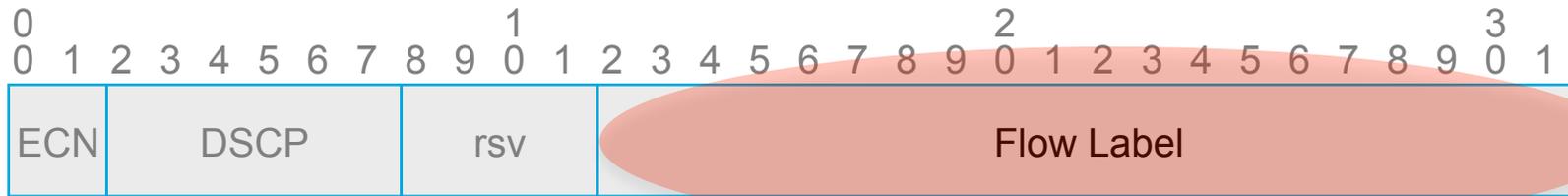


TF	2 bits	Traffic Class and Flow Label
NH	1 bit	Next Header
HLIM	2 bits	Hop Limit
CID	1 bit	Context Identifier Extension
SAC	1 bit	Source Address Context
SAM	2 bits	Source Address Mode
M	1 bit	Multicast Address Compression
DAC	1 bit	Destination Address Context
DAM	2 bits	Destination Address Mode

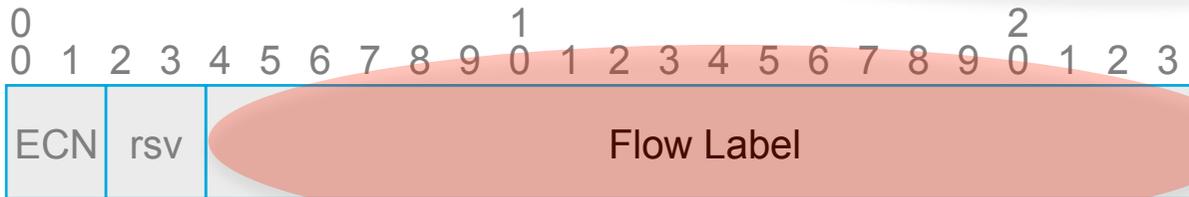
# SLoWPAN: Traffic Class & Flow Label



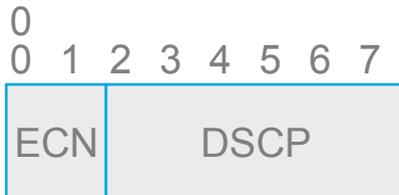
TF = 0



TF = 1



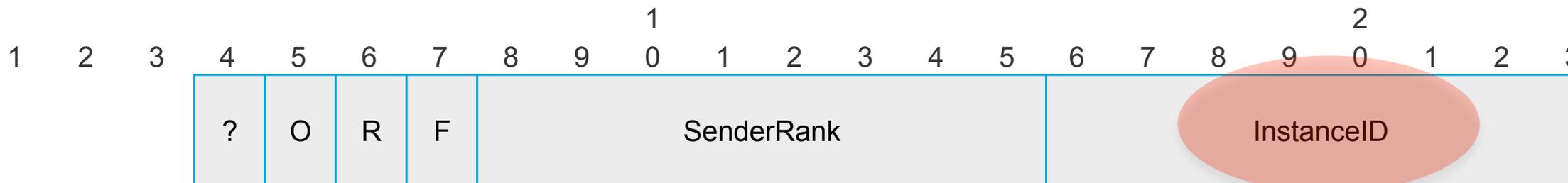
TF = 2



TF = 3

Traffic Class and Flow Label elided.

# draft-thubert-roll-flow-label



aces in Flow Label the RPL Packet Information is defined in RFC 6550 [Section 11.2](#)

ave extra HbH header bytes incurred in RFC 6553 AND eventual IPinIP tunneling

“When the router is the source of the original packet and the destination is known to be within the same RPL Instance, the router SHOULD include the RPL Option directly within the original packet. Otherwise, routers MUST use IPv6-in-IPv6 tunneling [\[RFC2473\]](#) and place the RPL Option in the tunnel header.”

discussed with Brian Carpenter on the ROLL ML

<http://www.ietf.org/mail-archive/web/roll/current/msg06967.html>

eed to take to 6MAN to resolve violation of RFC6437

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- ✓ 6TSCH
- ✓ Polymorphic Flows



# IIOT Network Convergence

Control anything from the network

Learn from Industrial.

Replicate and generalize with open standards

A converged network provides

high availability, flow isolation, security,

Guaranteed bandwidth and **Determinism**

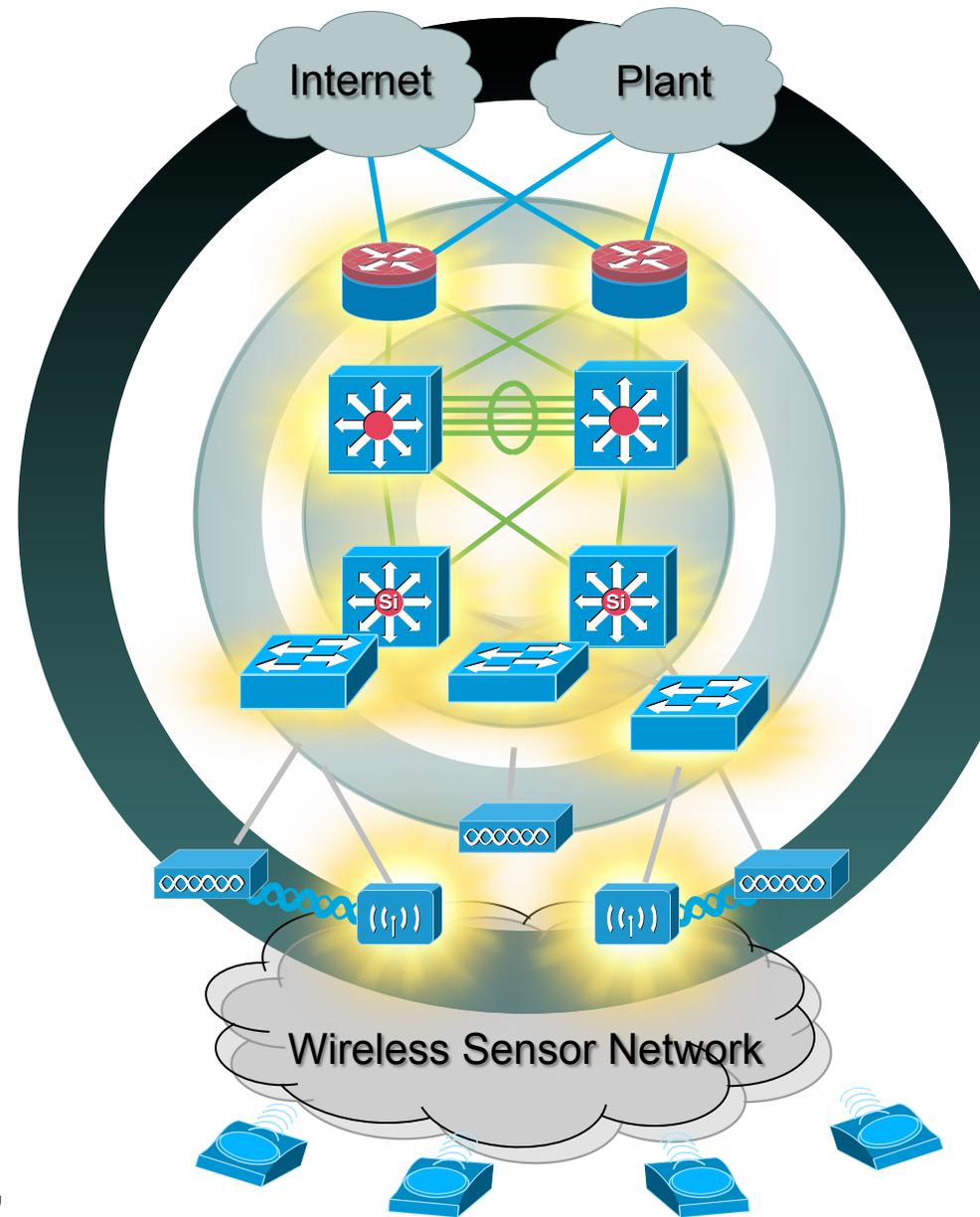
Coexistence involves:

Location and Scheduling

At L2 and L3, for wired and wireless.

New, higher end paradigm

Reaching more devices and data, farther, cheaper,  
with better guarantees



Thank you.

