

The Wireless FRINGE

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Advanced Architecture and Research

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Agenda

The Fringe of the Internet

The Fringe backbone

The Deterministic Fringe

6TSCH

Polymorphic Flows



The routing infrastructure, today



The Internet

Fully engineered

Hierarchical, Aggregations, ASs, Wire links

Fully distributed States

Shows limits (BGP tables, addr. depletion)

⇒ Reached adult size, mature to aging

⇒ Conceptually unchanged by IPv6

IPv4 Intranets

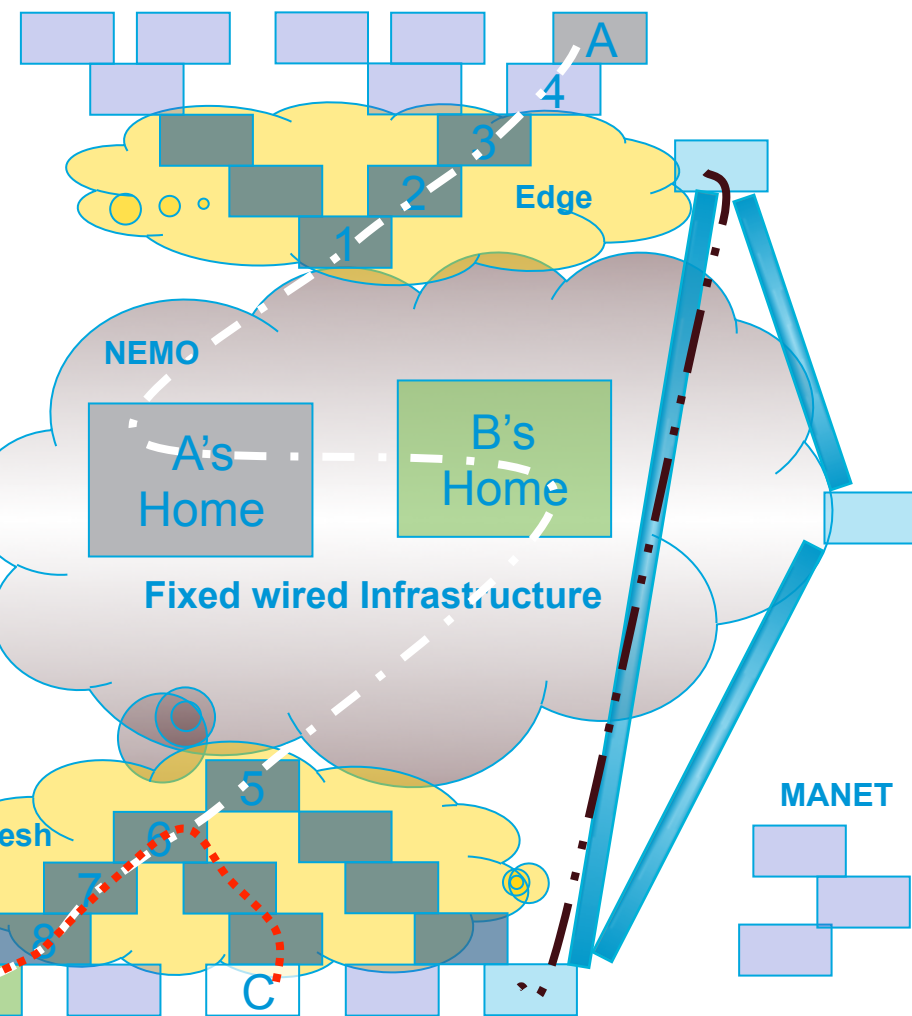
Same structure as the Internet

Yet decoupled from the Internet

NAT, Socks, Proxies

⇒ First model for Internet extension

The emerging Fringe of the Internet



L2 mesh Under

Multi-hop Public Access Points,
Proprietary mission specific products
Getting Deterministic

L3 Route Over

Migration to IETF Protocols (RPL)
Internet of Things (IOT, M2M)
Different IPv6 (6LoWPAN, SDN)

Mobile Overlays

Global reachability (NEMO, LISP, DMM)
DataCenter Network virtualization (NVO3, LISP)
Route Projection

**The Fringe DOES NOT LEAK
into the
Routing Infrastructure**

A sense of history

Router only knows “self” with: ID, certificates
Peers are discovered
Links are discovered
Routes are discovered
→ Infinity of self-centric networks

IPv6 autonomic

scale

IPv4 routing

Router CLI with: ID, keys.
All links to L2 peers
Routes are discovered
=> Single ‘GRID’

NA Subarea

NCP generation with:
All Transmission Groups to L2 peers
All Physical Units type 4 nodes,
All Virtual Routes

Scaling to Pervasive IoE/T

000*scale => No leak in the Internet

=> Opaque Fringe operations

=> Federating backbone

Reachability

=> Radio

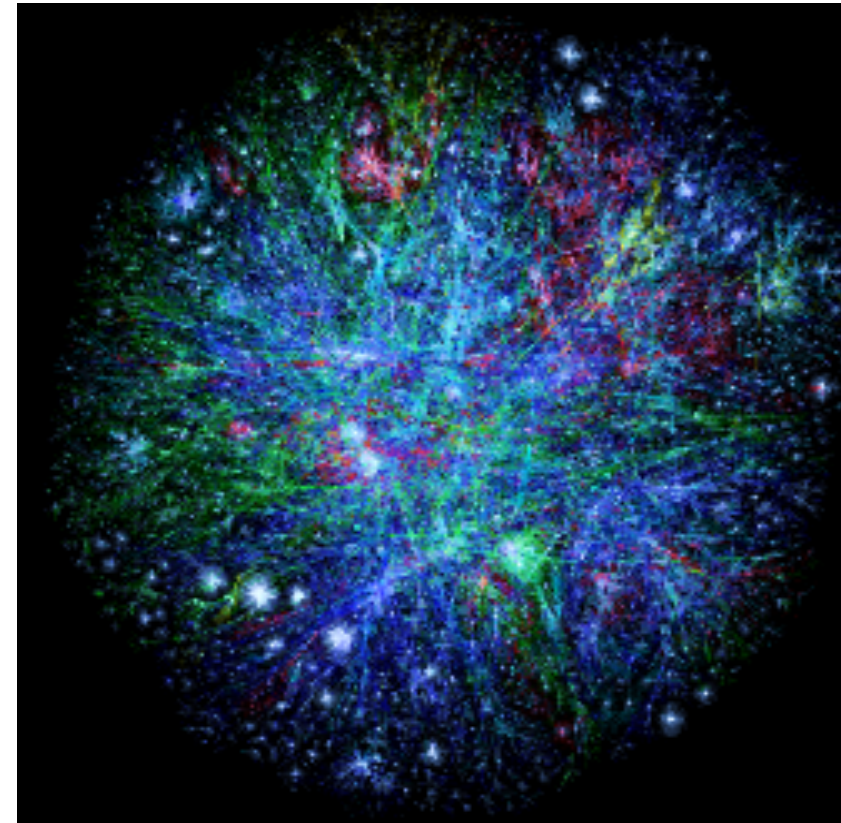
Addressing

=> IPv6

Density

=> spatial reuse

=> Routing



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Backbone Devices

Wired or Wireless

Switches and Controllers

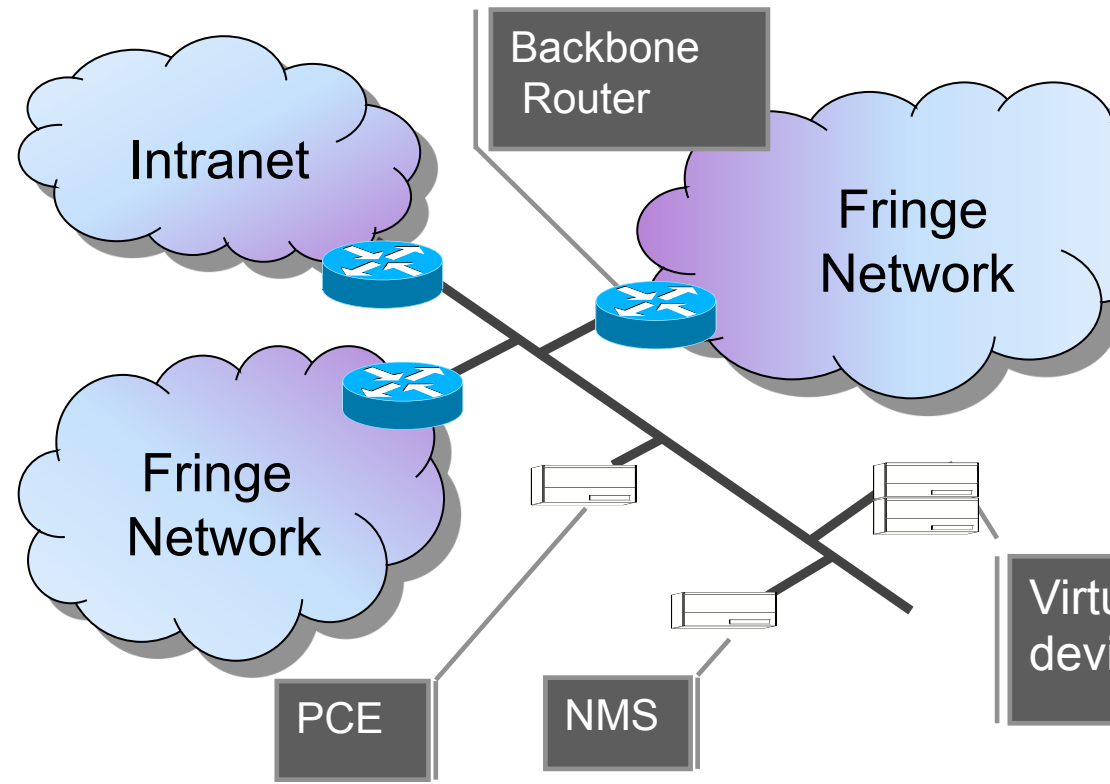
Wireless ND

Global Mobility integration

Virtualized Functions

PCE for deterministic route
computation

Legacy Appliances e.g. industrial PLC



The IPv6 Neighbor Discovery (R)evolution

A new Efficient ND, aka WiND for Wireless ND

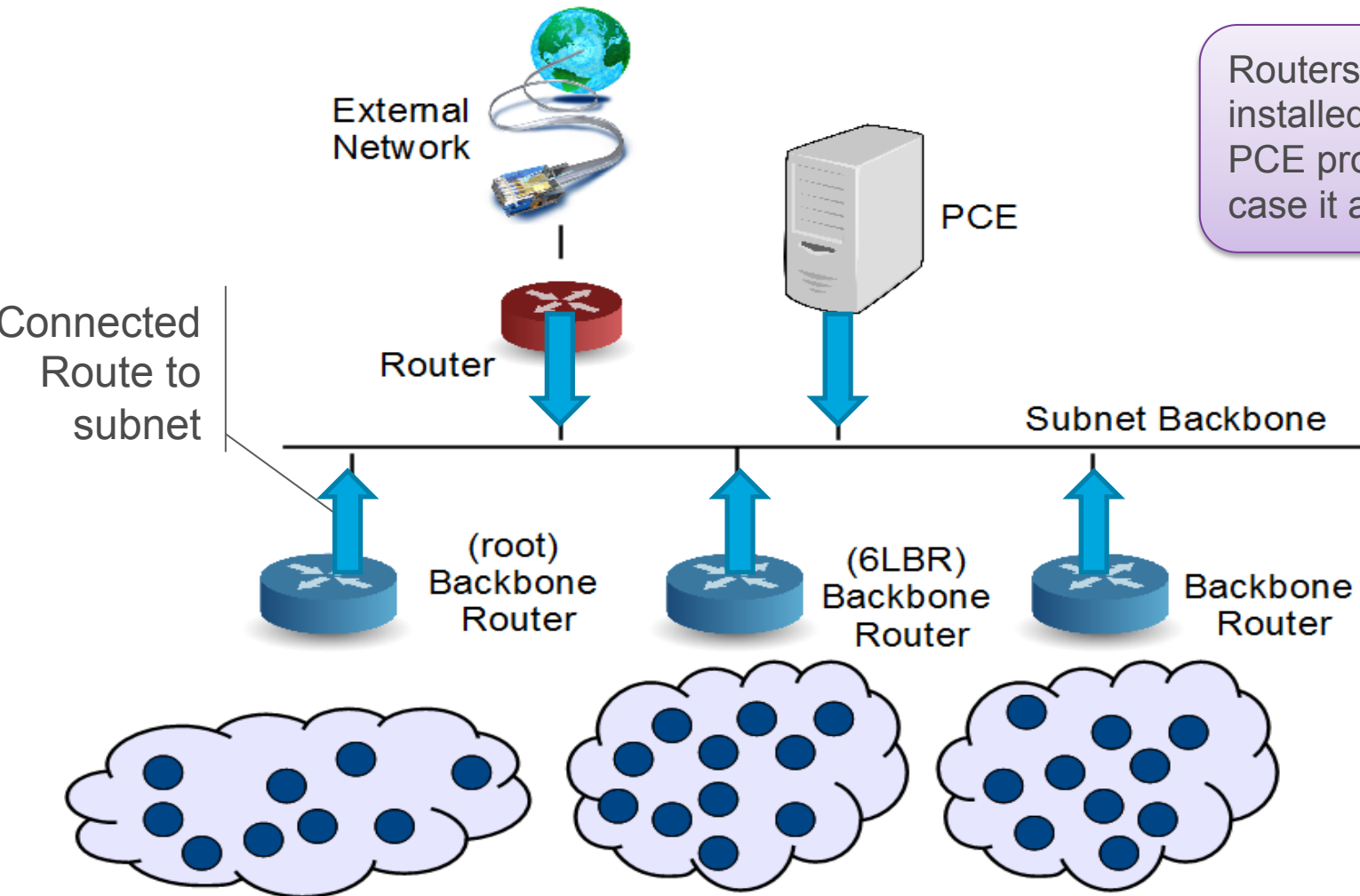
Registration and Duplicate Address Detection (DAD)
extending 6LoWPAN ND over the backbone

Resolution

Distributed through proxy-ND

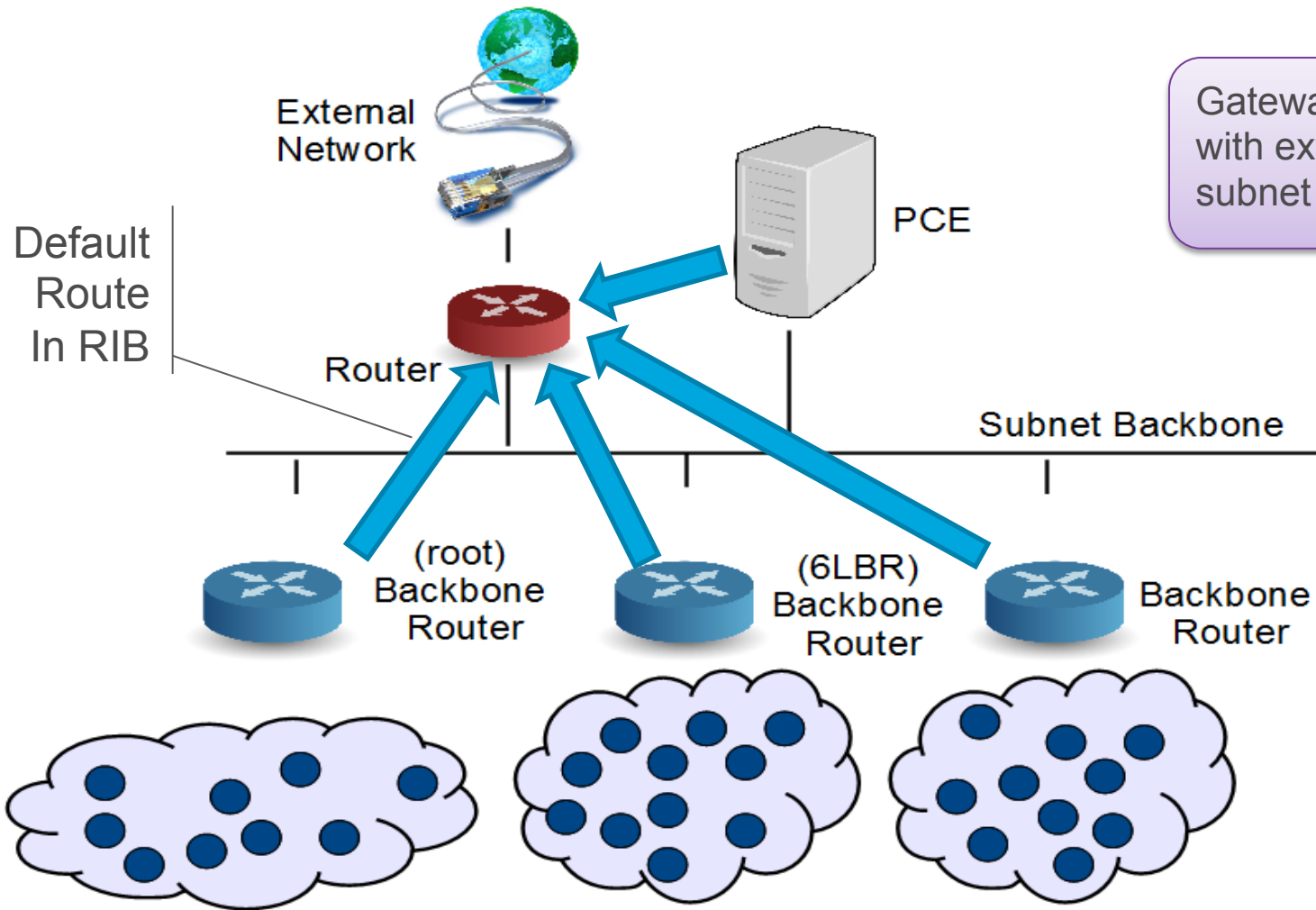
Routing in not-onlink mode

Initial time



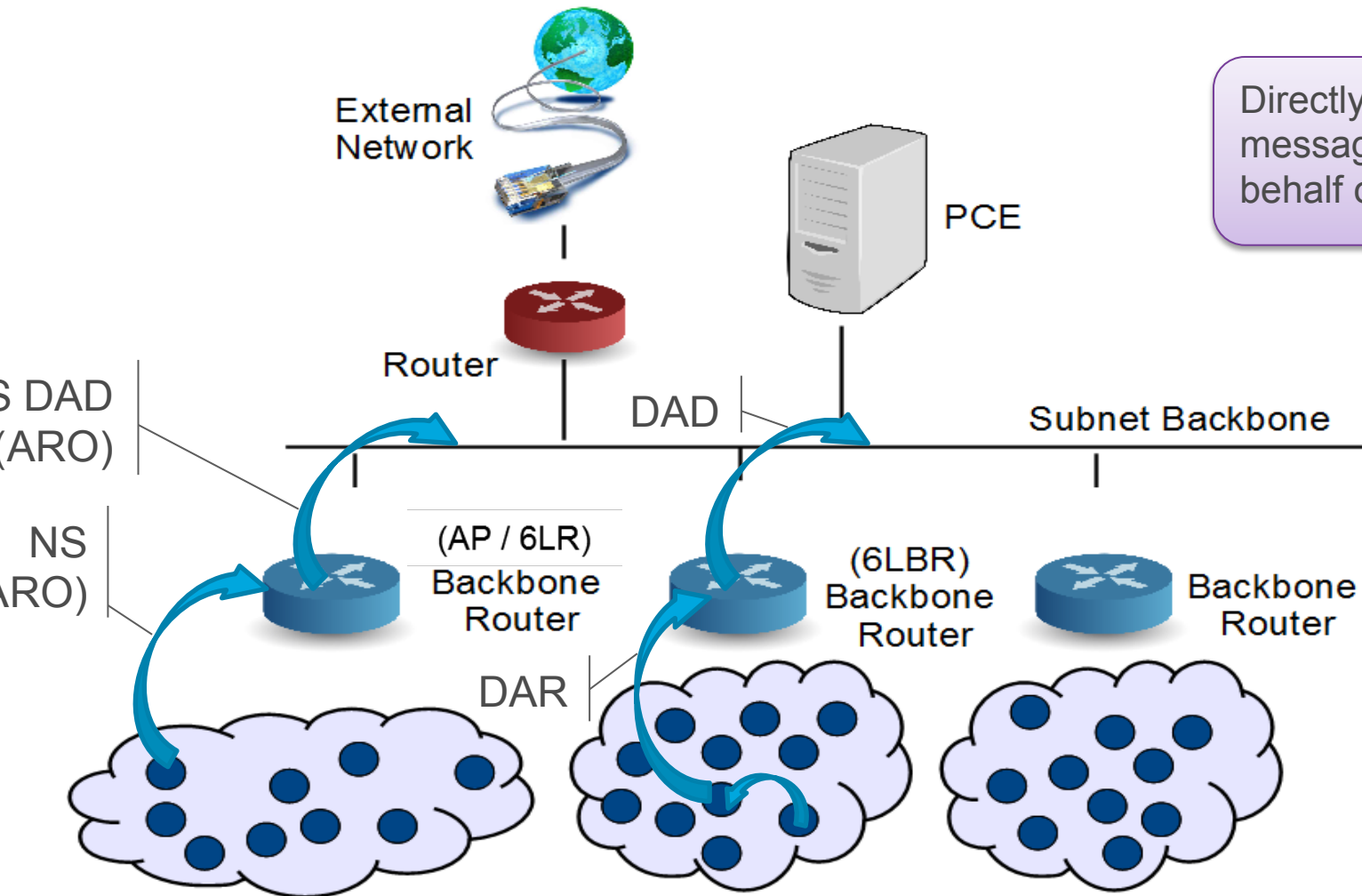
Routers within subnet have a connected route installed over the subnet backbone.
PCE probably has a static address in which case it also has a connected route

First advertisements from GW (RA, IGP, RPL)



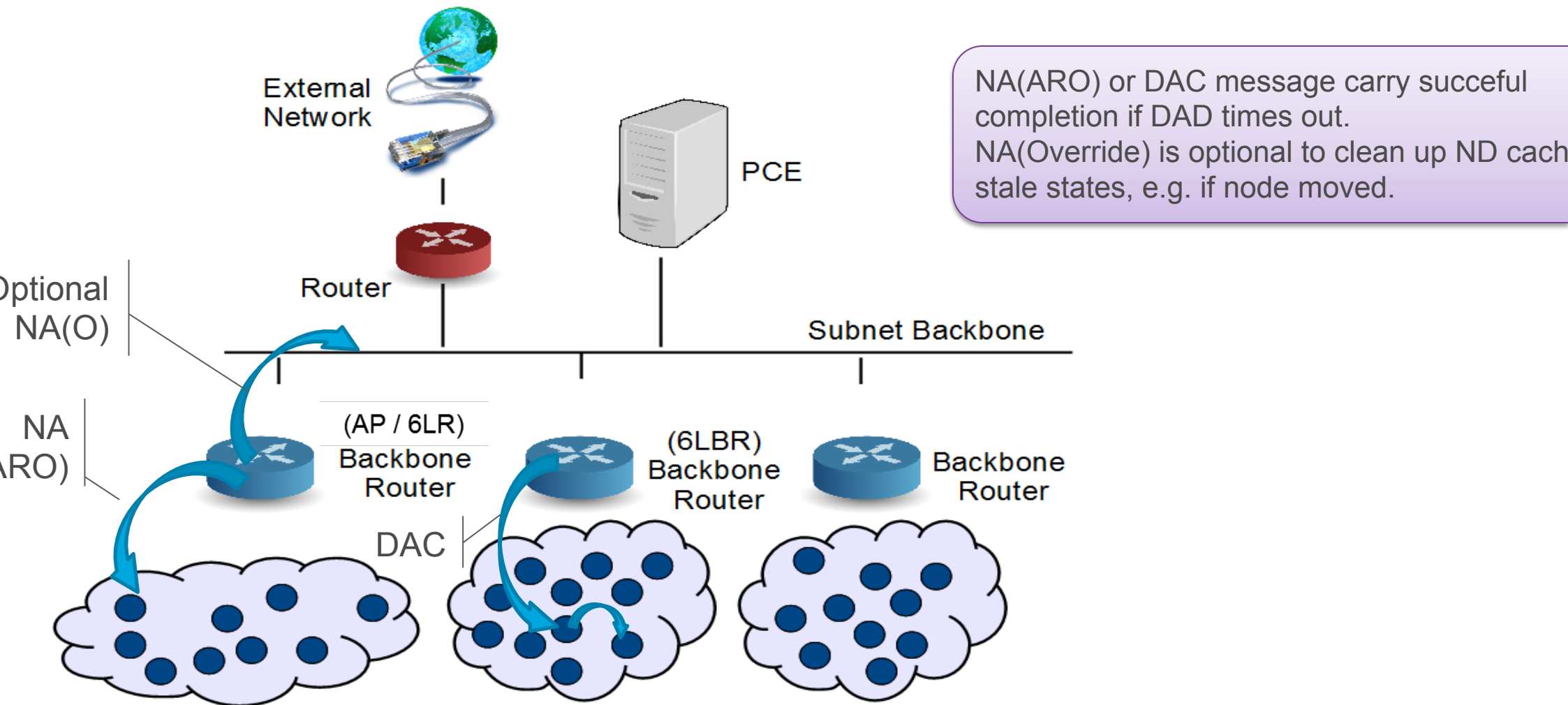
Gateway to the outside participate to some IGP with external network and attracts all extra-subnet traffic via protocols over the backbone

IPv6 ND Registration to 6LR and 6LBR

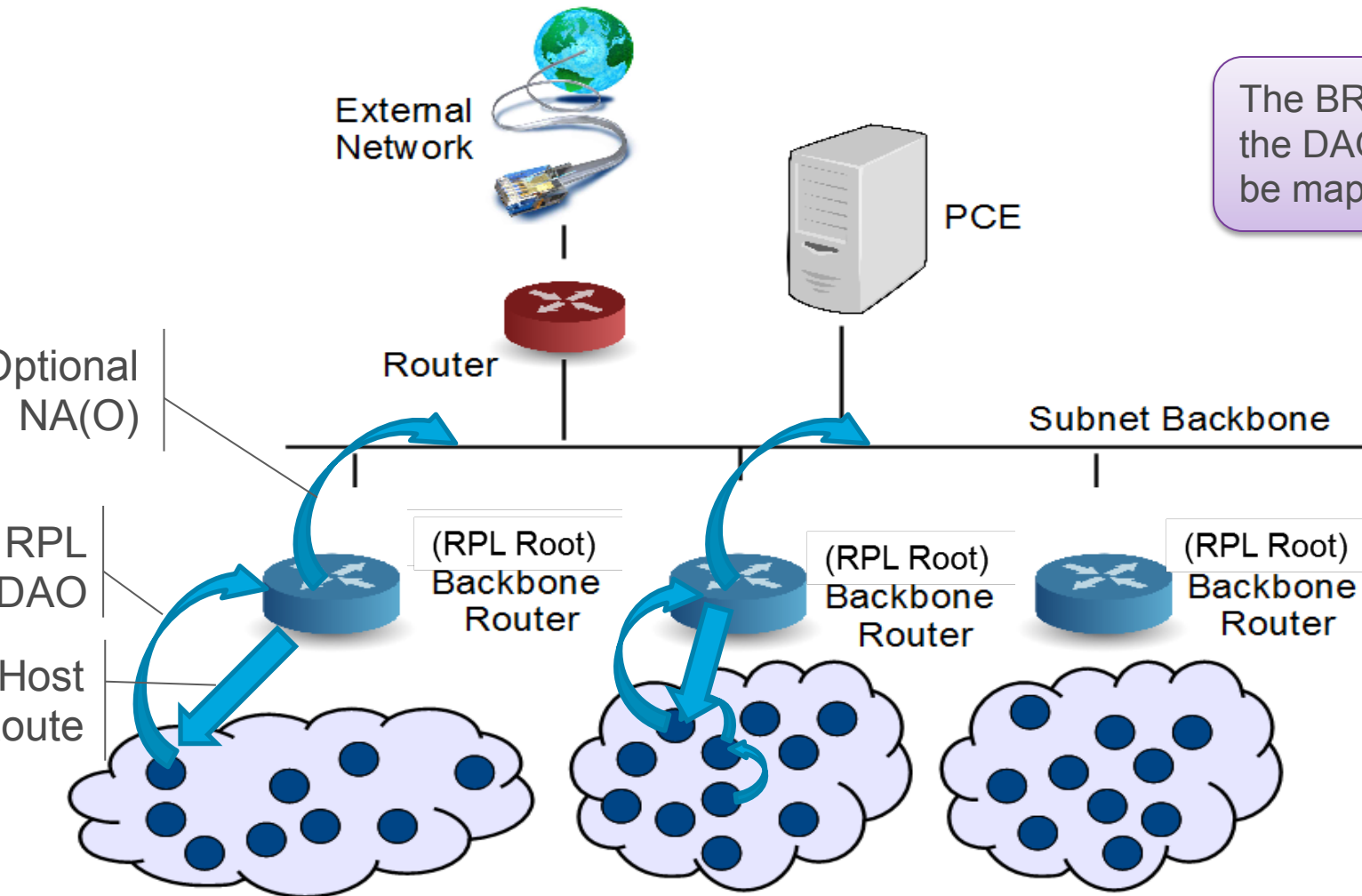


Directly upon NS(ARO) or indirectly upon DA message, the backbone router performs DAD behalf of the wireless device.

IPv6 ND Registration and Proxy for NS ARO

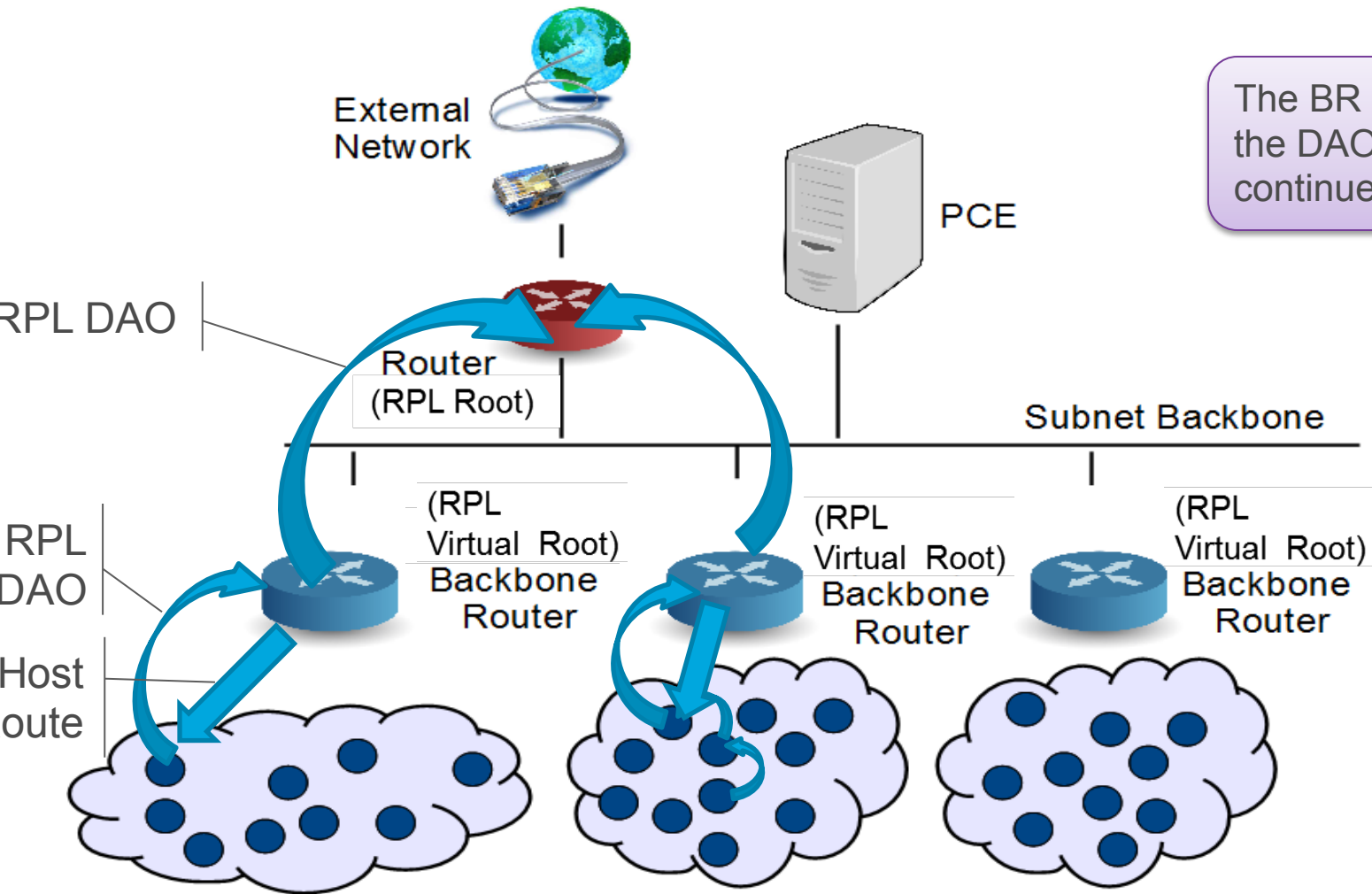


IPv6 ND Proxy for RPL



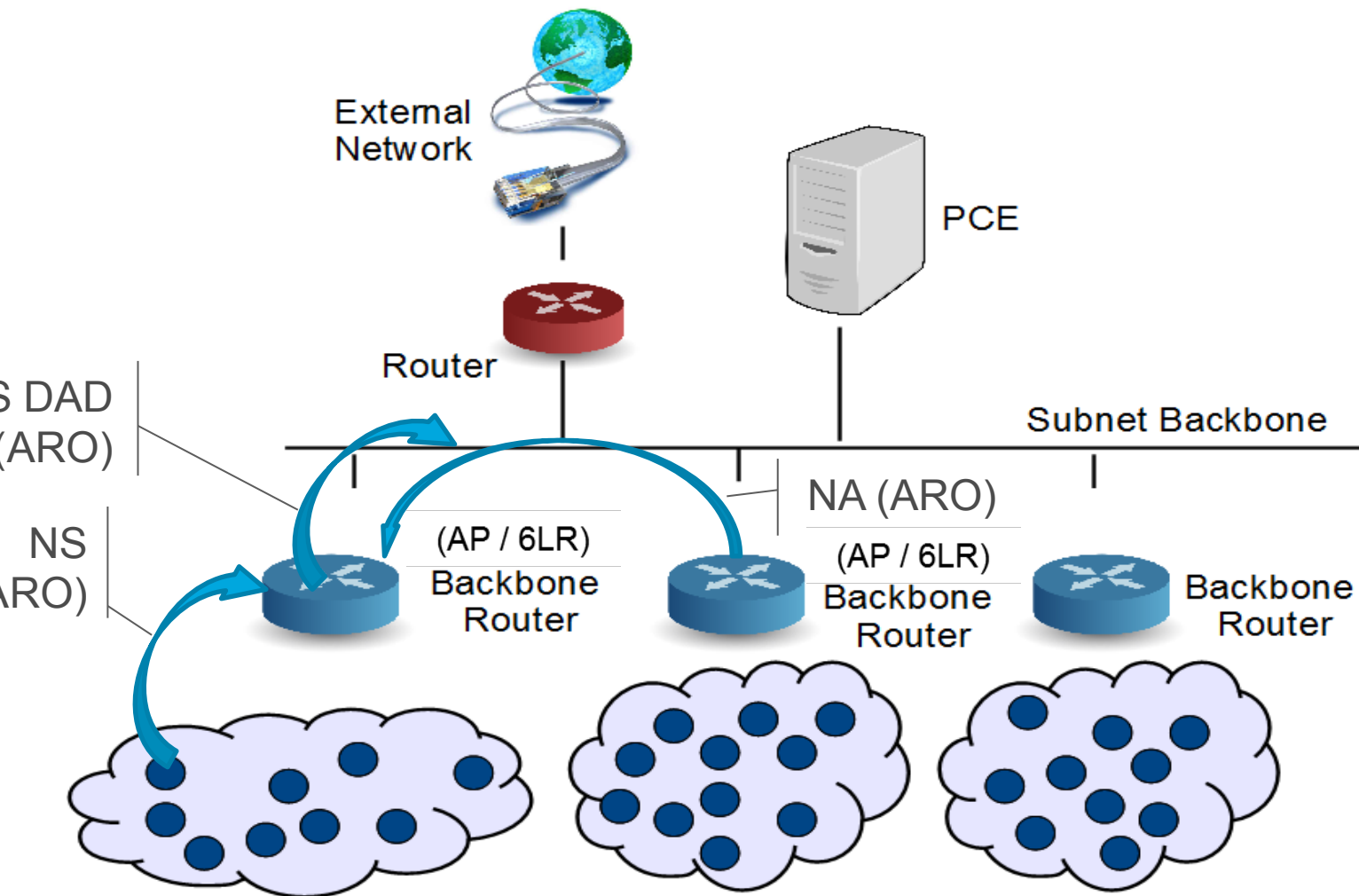
The BR maintains a route to the WSN node for the DAO Lifetime over instance VRF. VRF may be mapped onto a VLAN on the backbone.

RPL over the backbone



The BR maintains a route to the WSN node for the DAO Lifetime over instance VRF that is continued with RPL over backbone.

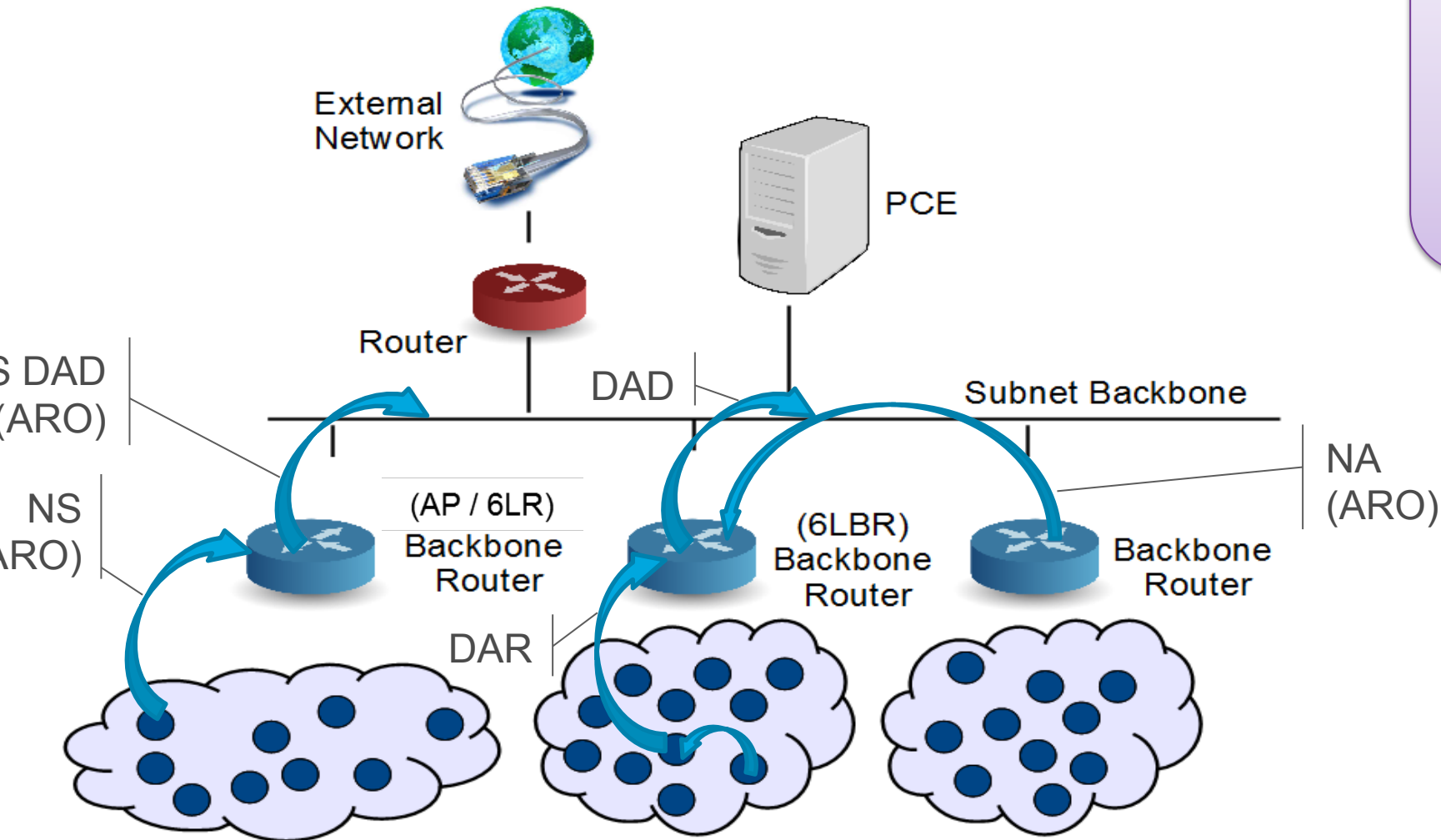
Duplication



DAD option has:
Unique ID
TID (SeqNum)

Defend with NA if:
Different OUID
Newer TID

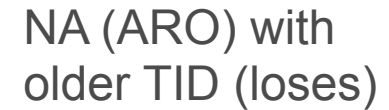
Mobility



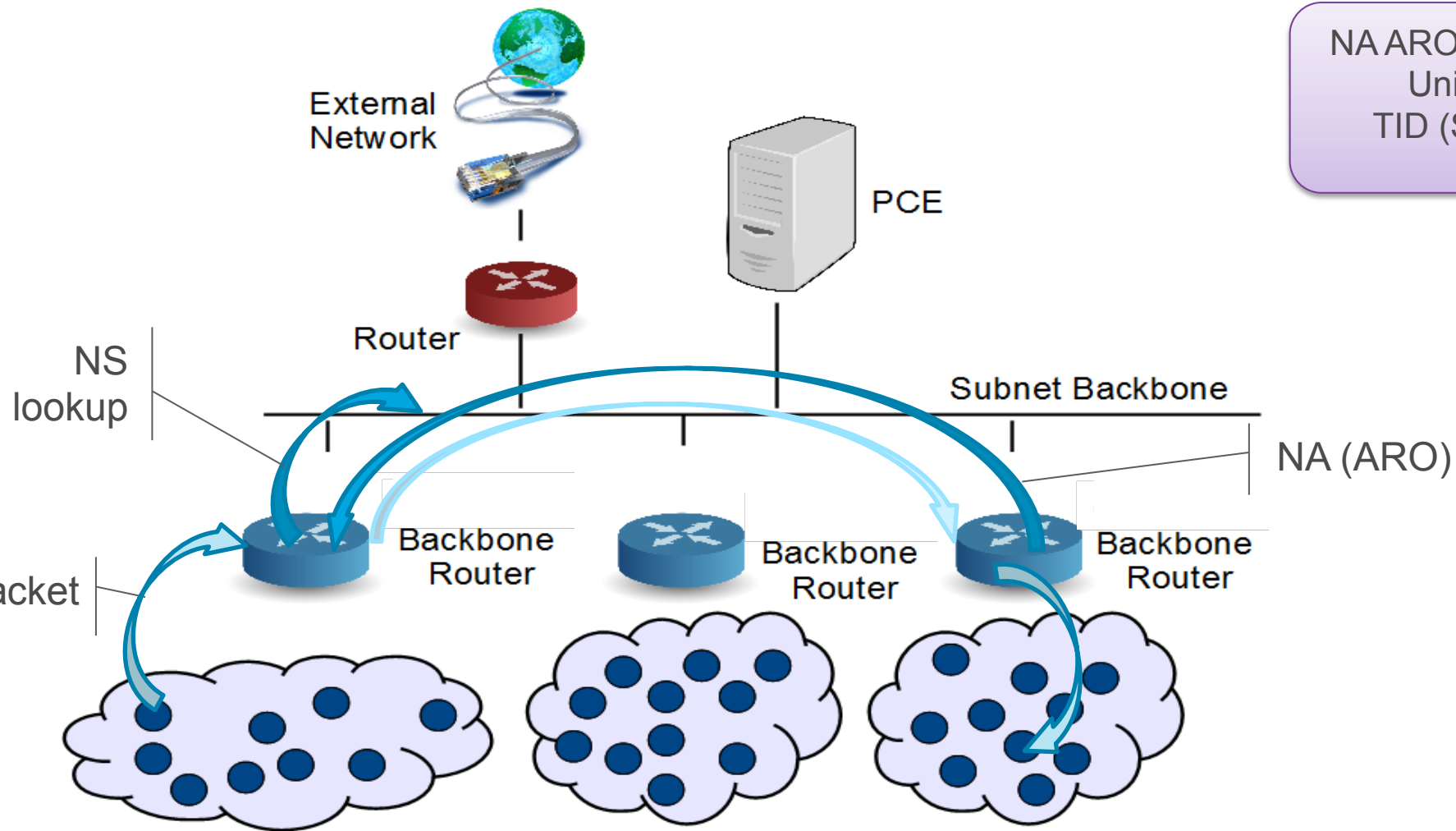
DAD option has:
Unique ID
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Defend with NA if:
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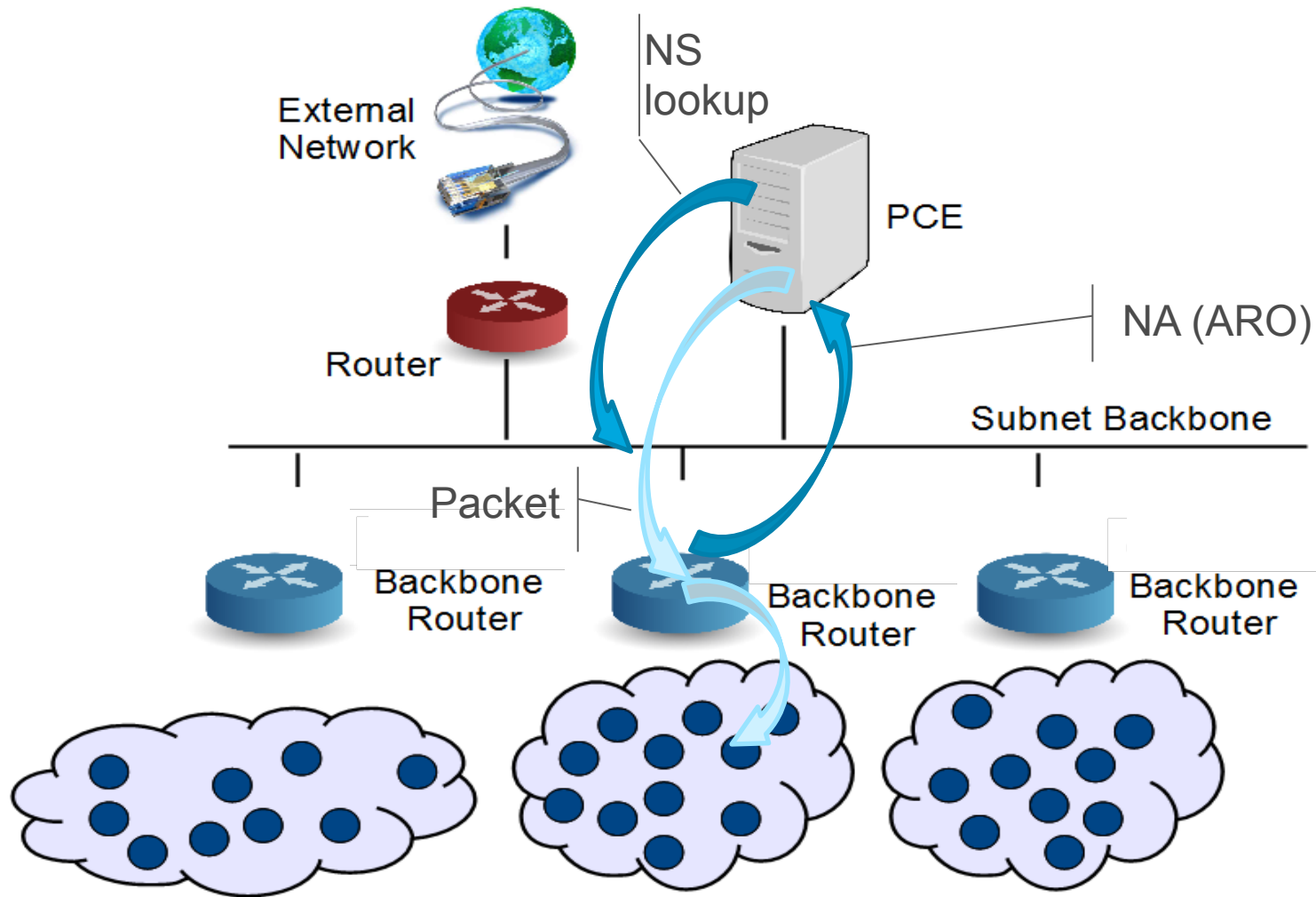
Host
oute



Resolution



Resolution (2)



Mixed mode ND
BBR proxying over
the backbone

Agenda

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- ✓ The Fringe backbone
- The Deterministic Fringe
- 6TSCH
- Polymorphic Flows



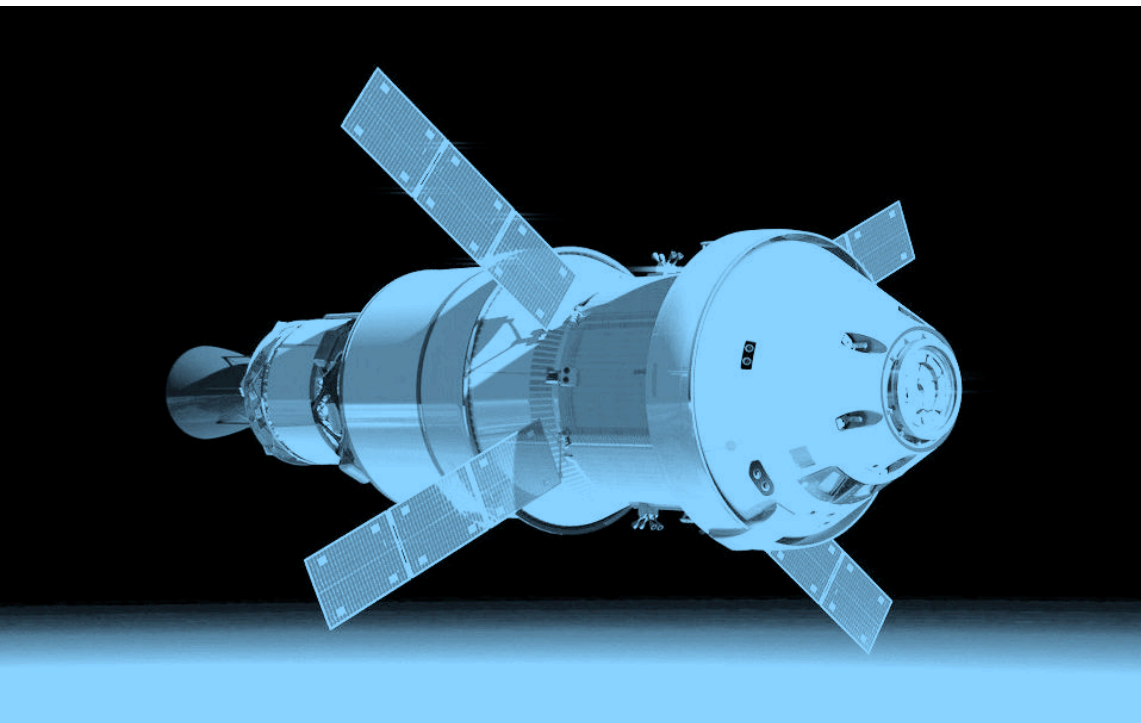
Deterministic Networking

new level of guarantees

For deterministic traffic (known a priori)

Network synchronization and Timely transmission

Centralized routing and scheduling optimization



Cleveland, Columbus & Cincinnati R.		
SPECIAL TIME SCHEDULE		
FOR THE TRAIN CONVEYING THE		
REMAINS OF ABRAHAM LINCOLN, LATE PRESIDENT OF THE U. S., AND		
FROM WASHINGTON, D. C., TO SPRINGFIELD, ILL.		
Cleveland to Columbus, Saturday, April 29th, 1865		
Leave Cleveland.....	12.00	Midnight.
Berea.....	12.43	A. M.
Olmsted.....	12.51	"
Columbia.....	1.02	"
Grafton.....	1.23	"
La Grange.....	1.37	"
Wellington.....	2.00	"
Rochester.....	2.17	"
New London.....	2.36	"
Greenwich.....	2.59	"
Shiloh.....	3.19	"
Shelby.....	3.39	"
Crestline.....	4.07	"
Galion.....	4.23	"
Iberia.....	4.41	"
Gilead.....	5.05	"
Cardington.....	5.20	"
Ashley.....	5.43	"
Eden.....	5.55	"
Berlin.....	6.19	"
Lewis Centre.....	6.32	"
Orange.....	6.47	"
Worthington.....	6.58	"
Arrive Columbus.....	7.30	A. M.
This Train will have exclusive right to the Road against all other		
Passenger Trains. The Locomotive will be run ten minutes in advance of the		
regular time.		
E. S. FLINT, Superintendent		

Process Control example: a refinery

Sensors and actuator usually close: <100m

Elements of critical loops can be wire-interconnected

Low-quality (noisy, intermittent) field power often available

Control room typically 500m to 2km distant

Wired: power + signal carried together on one twisted pair

Designed to meet intrinsic safety regulations: ~40mW/pair

Typically 25-pair to 100-pair wire bundles in buried conduit

Wireless is most attractive for this long, costly link

Primary requirement is ≥ 5 yr battery life for field devices

Battery replacement is often very costly or impractical

Environmental power harvesting strongly desired



e.g. ExxonMobil Baytown refinery near Houston (Texas)

100 hectares (40 sq. miles) with 12 control rooms (CRs): 10 refining CRs localized to a few hectares each,

2 CRs that span the plant and interact with the other 10; large chemical plant adjacent to refinery

What's specific

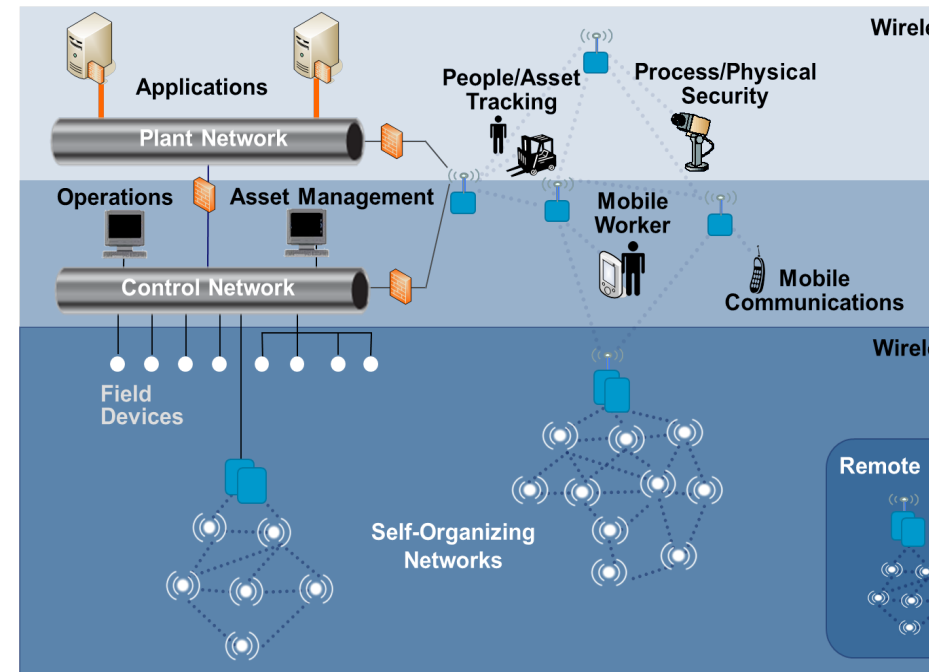
Industrial networks evolved concurrent with IP

Process Control and Factory Automation goals
differ from those of IT => schism

Need low latency and high cyclic determinism

Need higher reliability, faster repair, lesser maintenance

Largest plants potentially reap the most profit from small productivity improvements
thus more willing to innovate (in stages) and to fund proven new technology



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 - ✓ The Deterministic challenge
- 6TSCH
- Polymorphic Flows



6TSCH: standardizing IPv6 over TSCH MAC

Align existing standards

(RPL, 6LoWPAN, PANA, RSVP, PCEP, MPLS) over 802.15.4e TSCH

Support Mix of centralized and distributed deterministic routing

Design 6TUS sublayer for L3 interactions

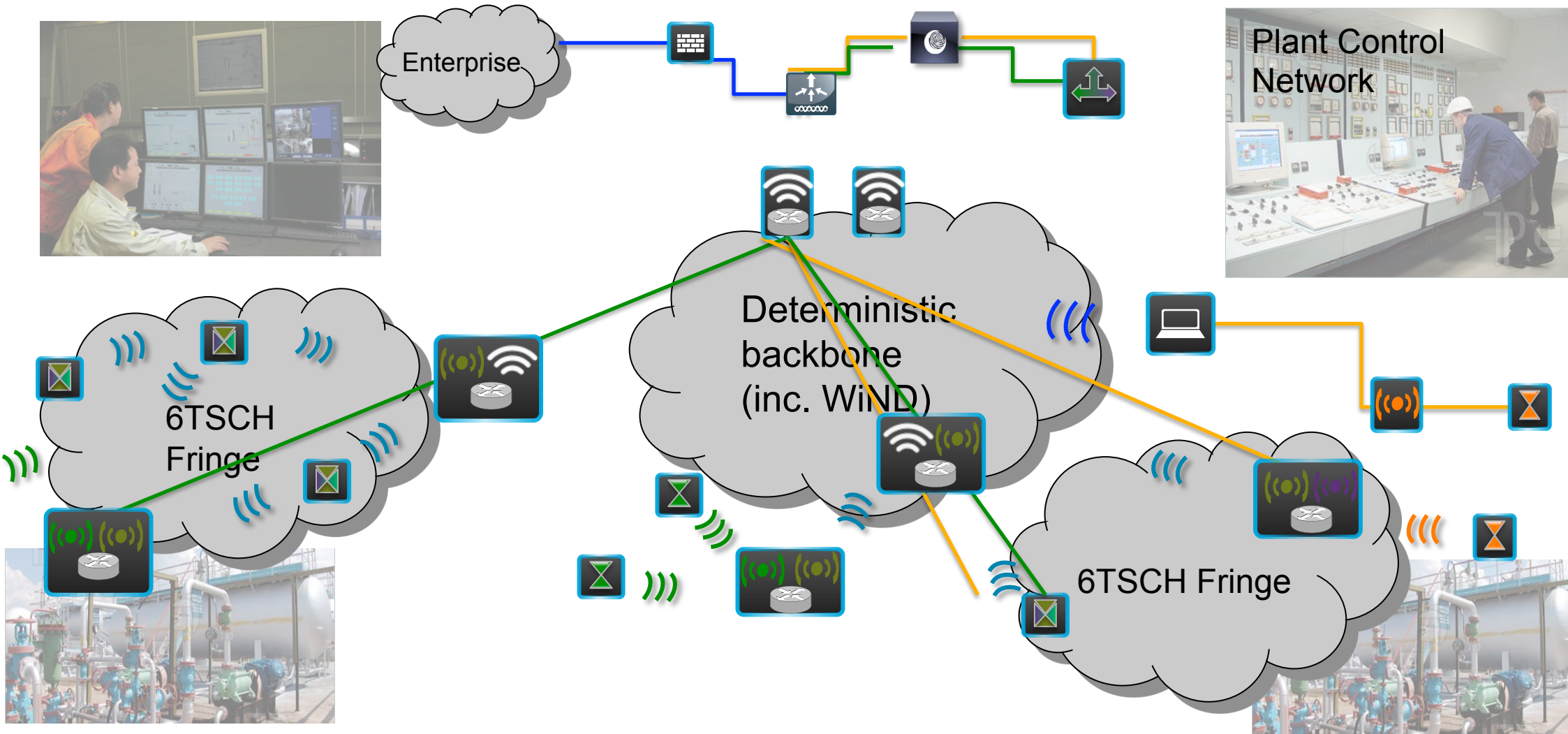
Open source implementations (openWSN...)

Build Ecosystem (DUST, NIVIS, IoT6, uni.lux, uni.bari, Berkeley...)

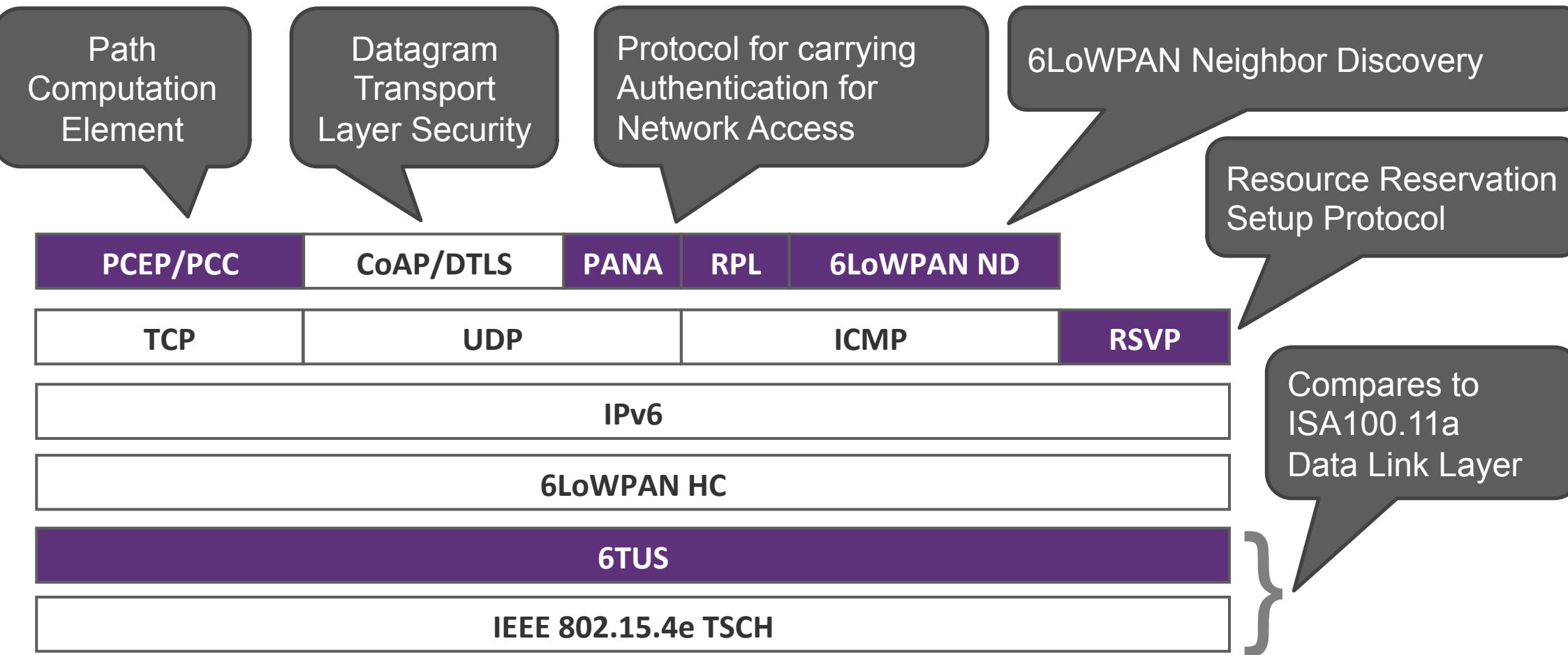
6 active drafts (Archi, 6TUS, ...)

Preparing for BoF in Berlin this summer

Future Industrial Architecture with unified wireless



TSCH: Architecture



Centralized vs. Distributed routing

Centralized

God's view optimization

Multipath redundancy

Deterministic (optimized)

Virtualization

Distributed

Autonomic & Mobile

Highly available (DARPA)

Deterministic

Scalability

Routing With RPL

Low Power Lossy Nets

Dynamic Topologies

Peer selection

Constrained Objects

Fuzzy Links

Routing, local Mobility

Global Mobility

Addressed in RPL ?



Distance Vector + stretch



Peer only with parents



DV + Non-storing mode



Lazy Update & datapath validation

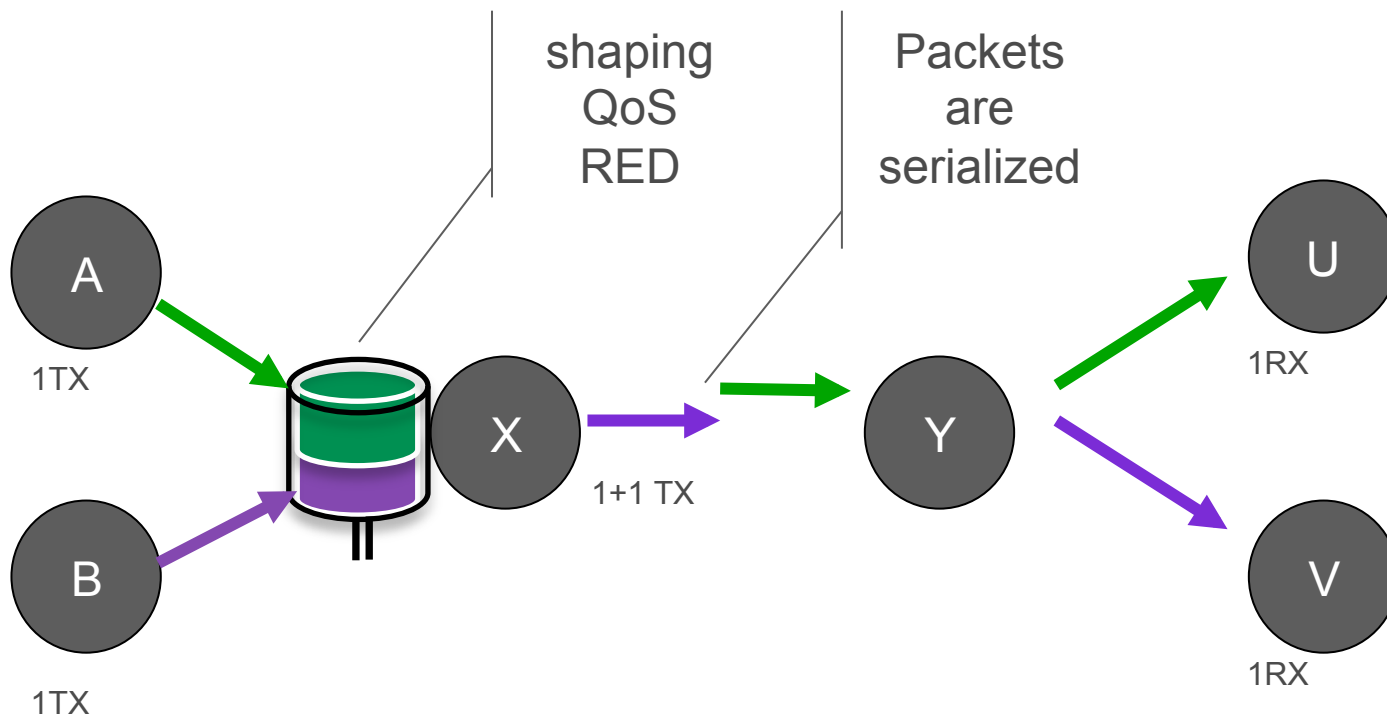


Constrained instances, TID

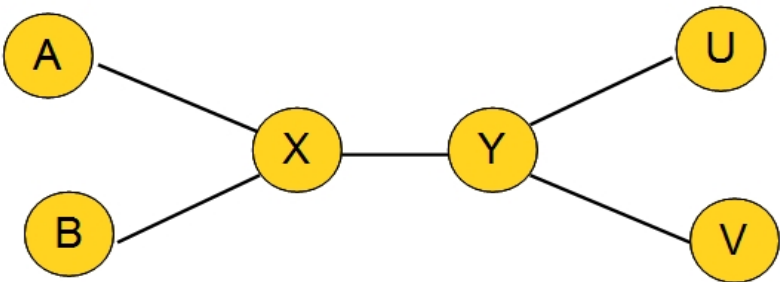


Req coupling with LISP/NEMO

Normal L3 operation



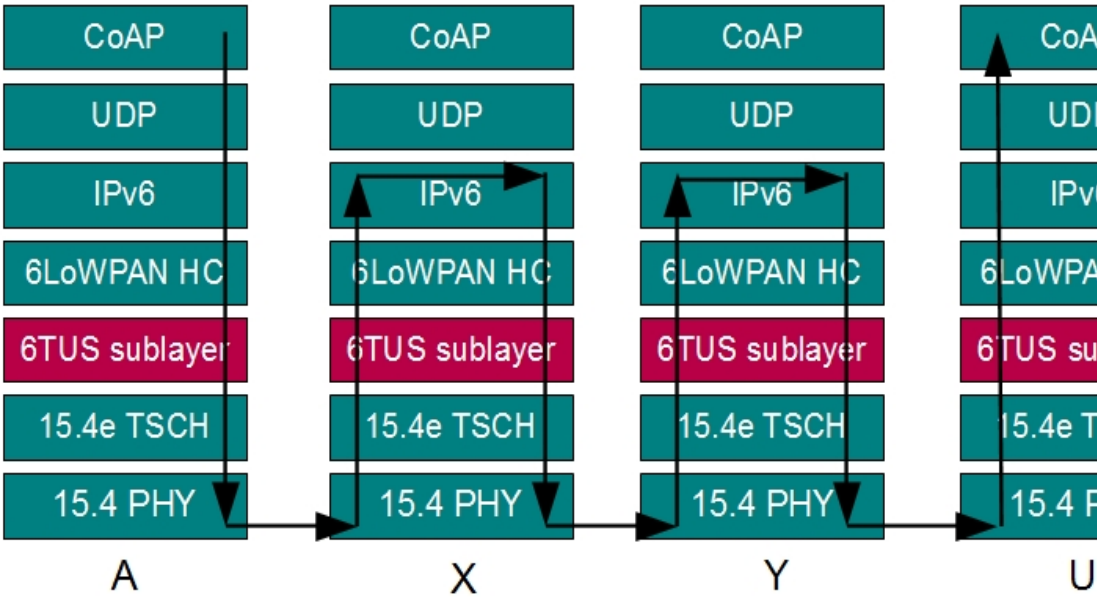
Best effort routing



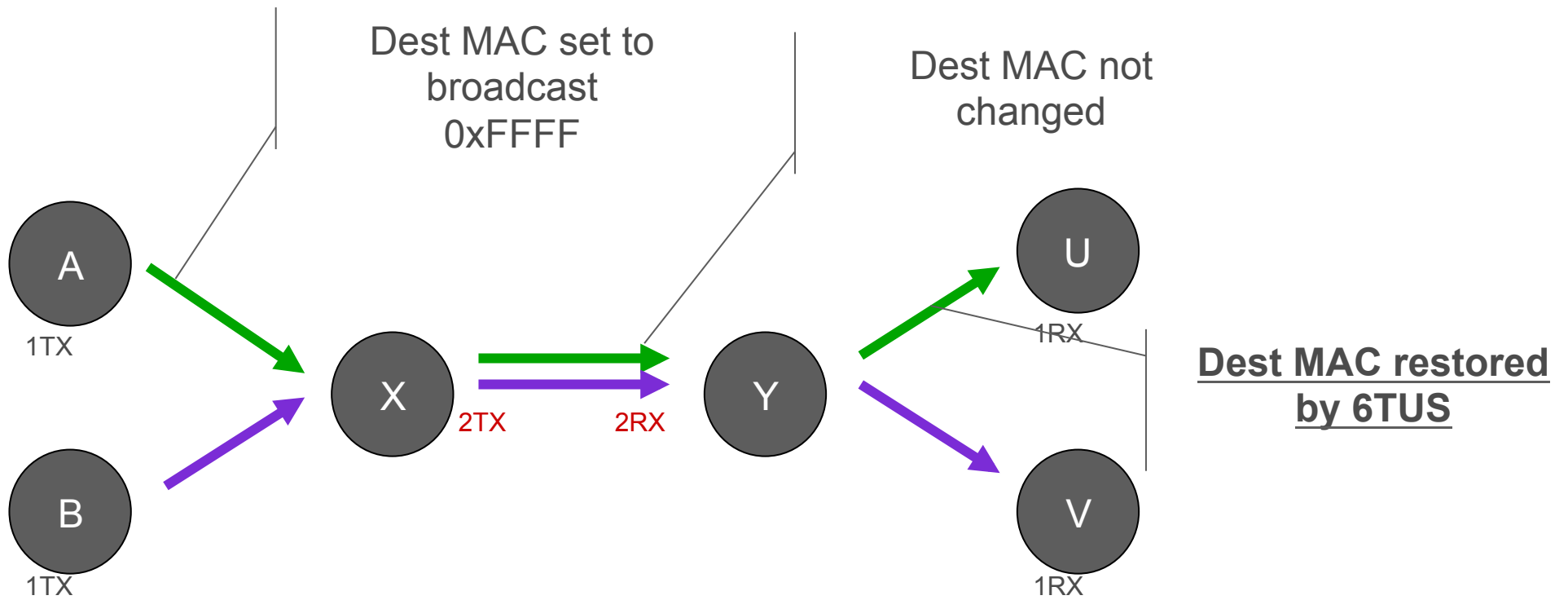
Bundle

			X → Y			Y → V	
A → X					X → Y		Y → U
	B → X						

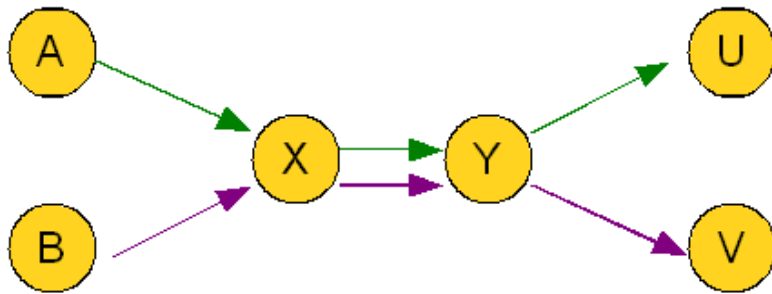
slotOffset



Normal Track operation



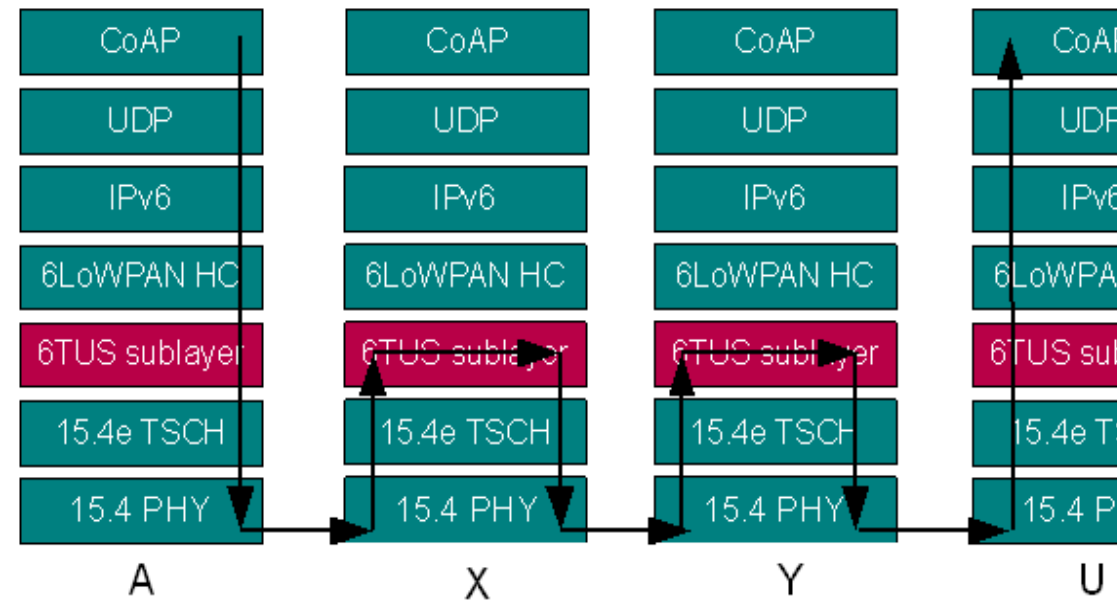
Track Switching (G-MPLS)



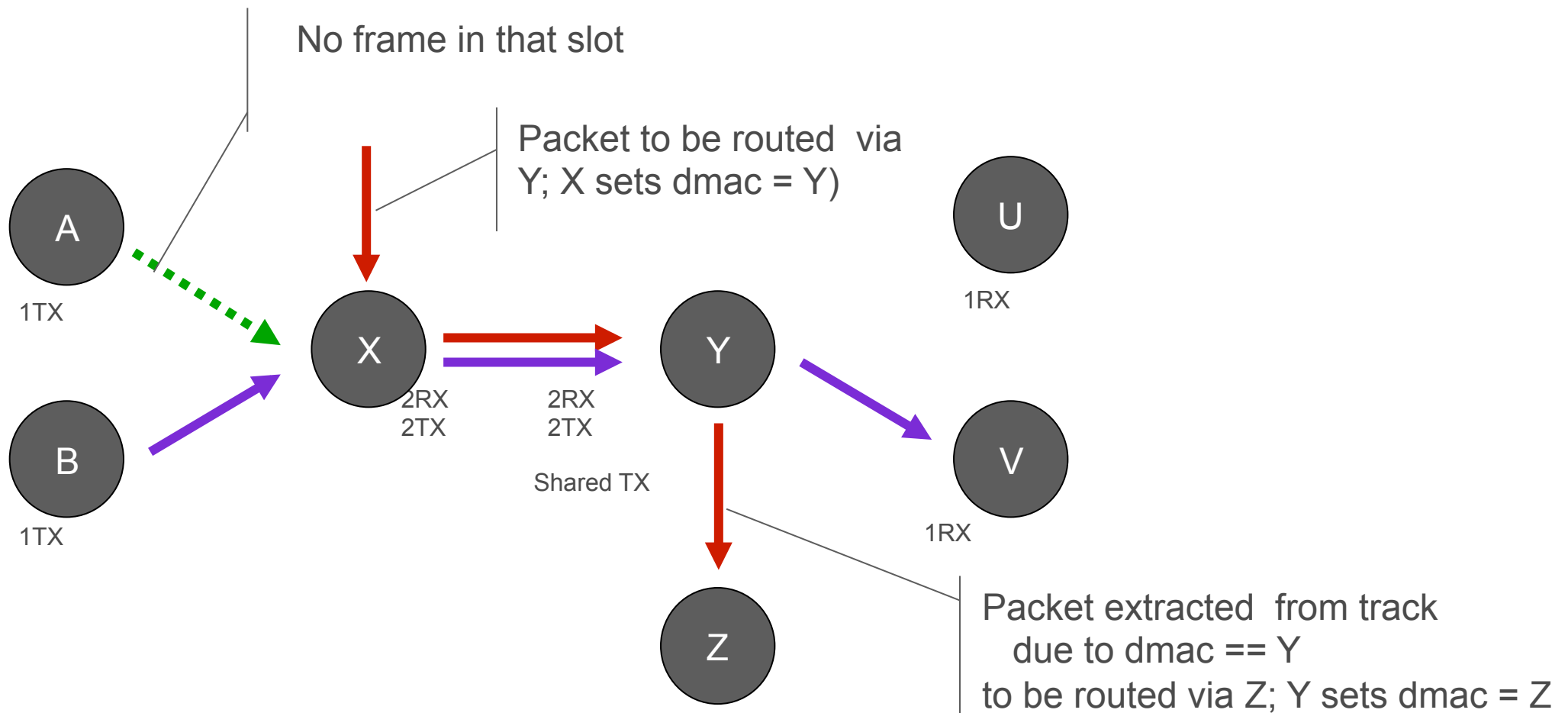
Track

			X → Y			Y → V	
		X → Y		X → Y			
A → X					X → Y		Y → U
	B → X						

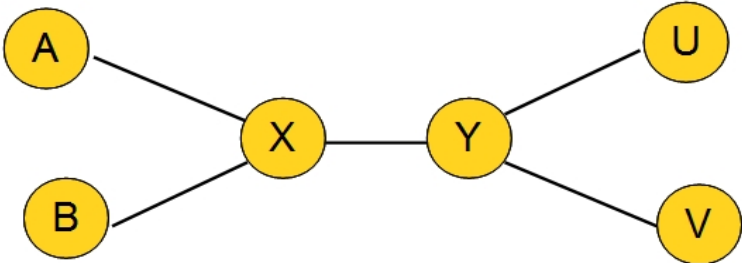
slotOffset



Opportunistic track slot reuse



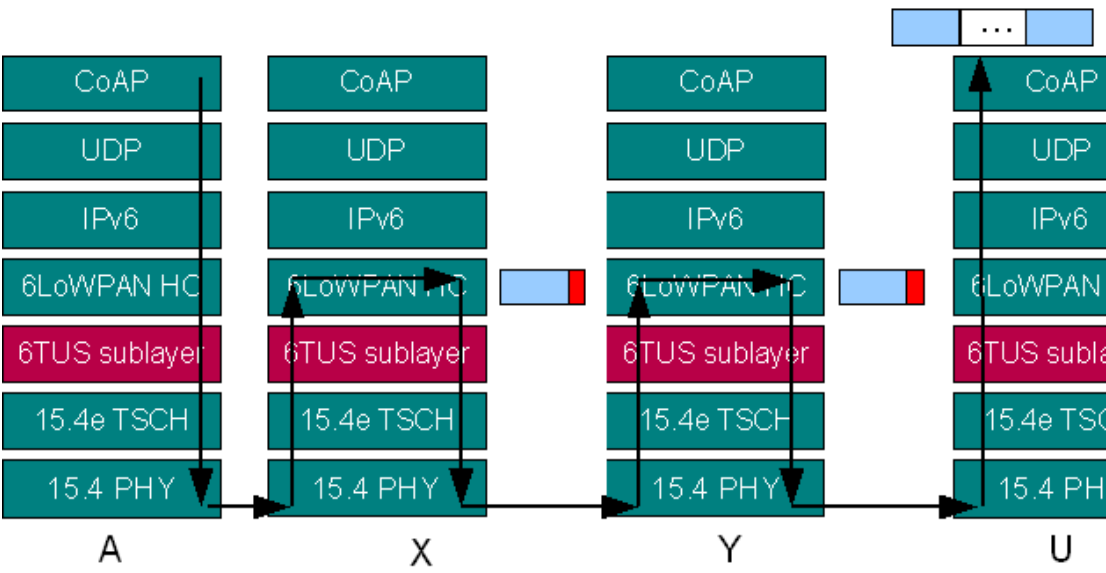
6LoWPAN Fragment forwarding



Bundle

			X → Y		Y → V	
A → X				X → Y		Y → U
	B → X					

slotOffset



packet, composed by N fragments
i-th fragment of the packet, with $i > 1$
state installed in a single fragment

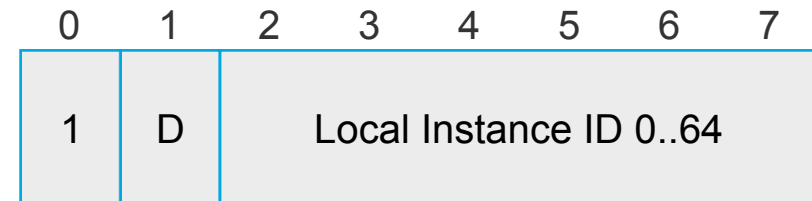
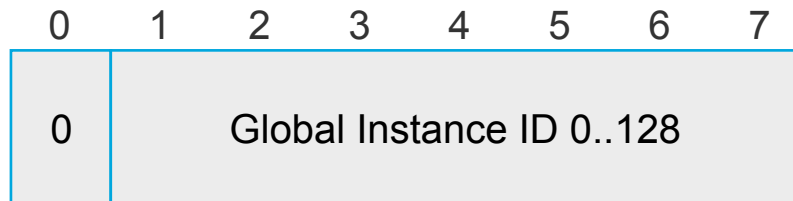
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RFC 6550: RPL Instance ID

The RPL instance ID allows different routing optimizations, constraints and policies.



The RPL instance ID is encoded in 1 octet. The first bit indicates whether Global or Local.

“A local RPLInstanceID is autoconfigured by the node that owns the DODAGID and it MUST be unique for that DODAGID. The DODAGID used to configure the local RPLInstanceID MUST be a reachable IPv6 address of the node, and it MUST be used as an endpoint of all communications within that Local instance.”

Inside a packet: “If the 'D' flag is set to 1, then the destination address of the IPv6 packet MUST be the DODAGID. If the 'D' flag is cleared, then the source address of the IPv6 packet MUST be the DODAGID.”

6TSCH extends RPL's language of DODAGID to route (reservation) endpoint.

Global Instances

28 global instances per network

Indexed by tuple (IPv6, InstanceID)

Running as Ships-in-the-night

1 instance = 1 VRF = 1 « L3 vlan »

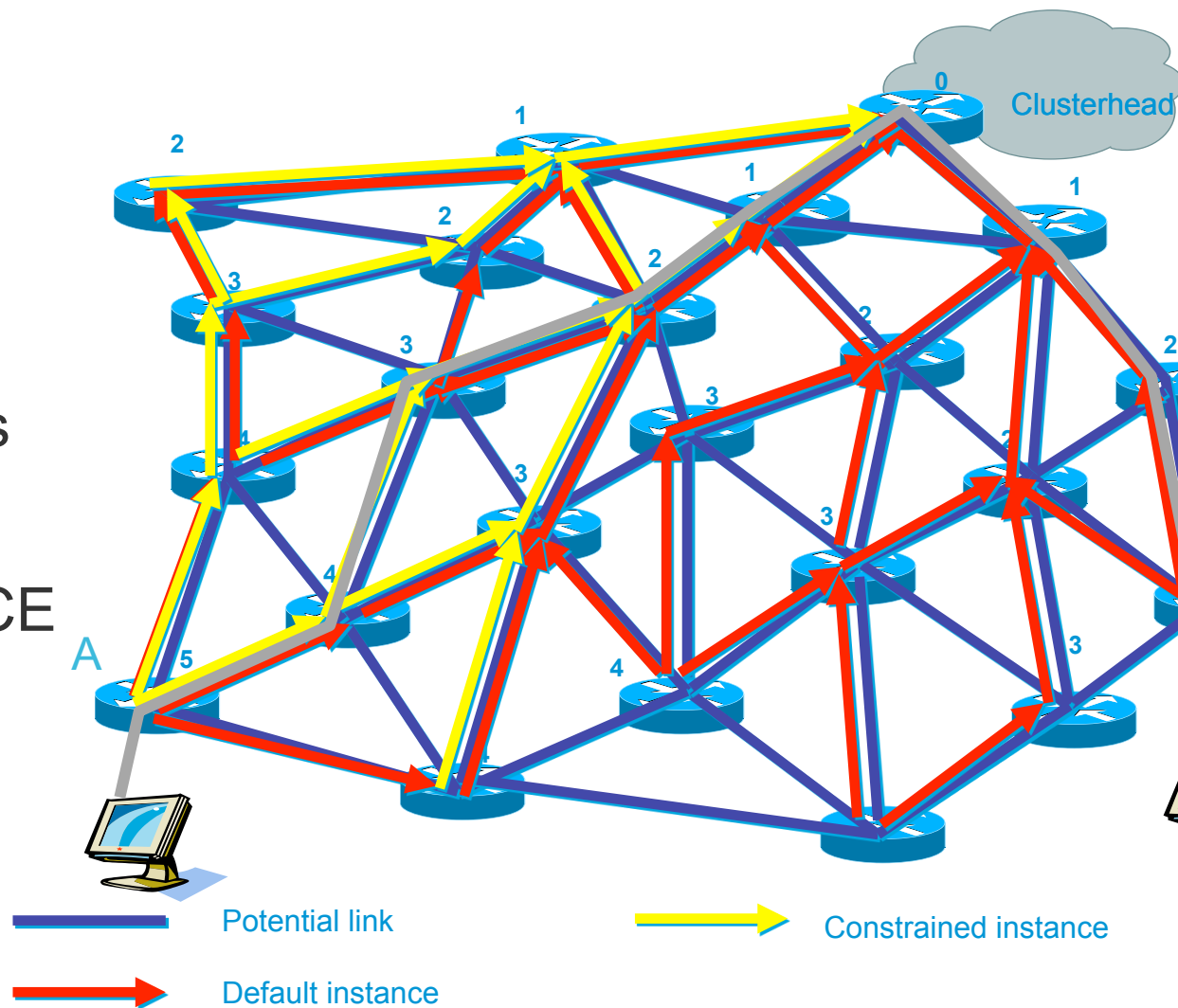
Serving different Objective Functions

Using different metrics

Can be shared between RPL and PCE

RPL along a DODAG

PCE adding orthogonal shortcuts



Local instances

64 local Instances

per IPv6 source address

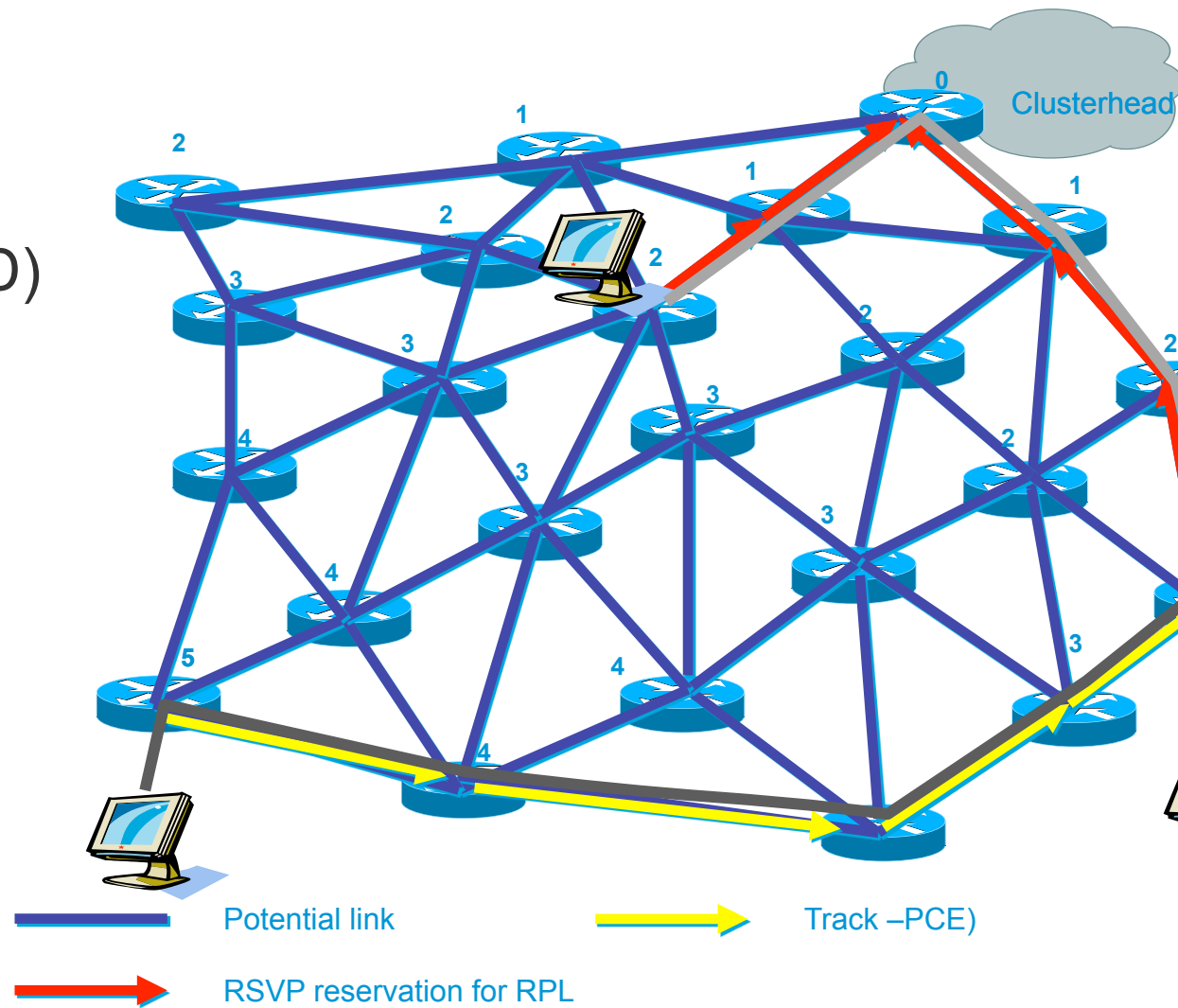
Indexed by tuple (IPv6, InstanceID)

Used by RPL or PCE

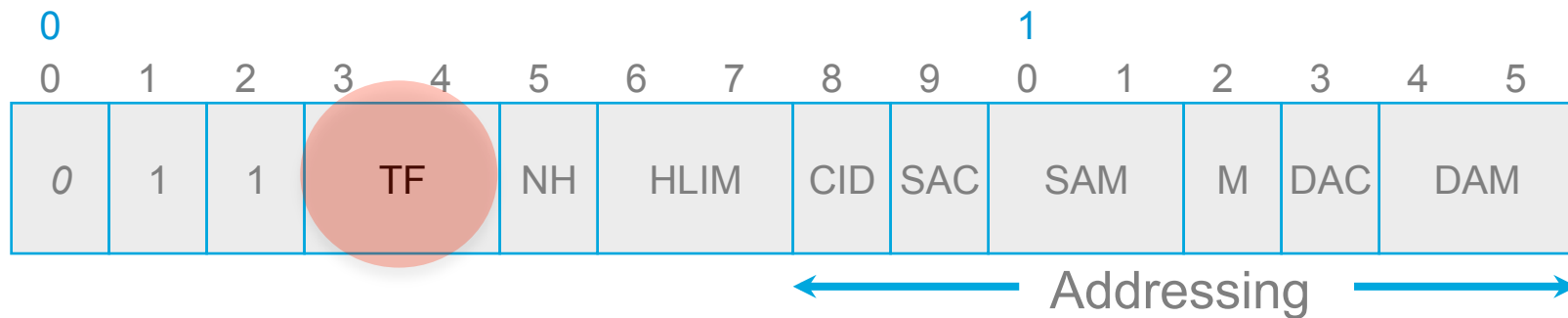
RPL: for P2P applications

RPL: to index RSVP path

PCE: Serves as Track ID,
included in PCEP request
from the source device

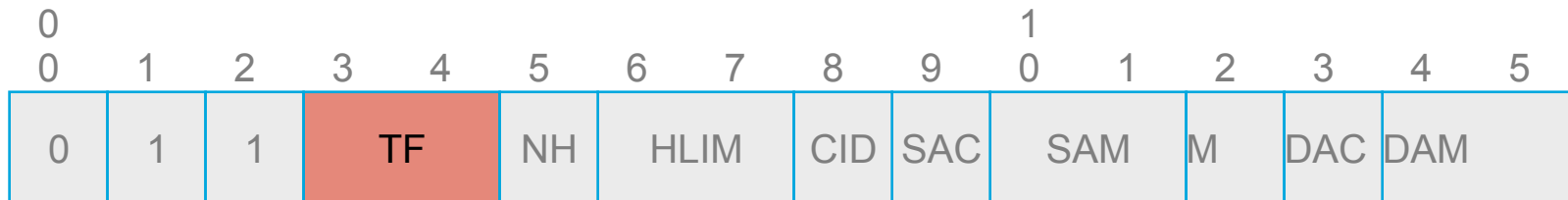


RFC 6282: 6LoWPAN IPv6 Header Compression

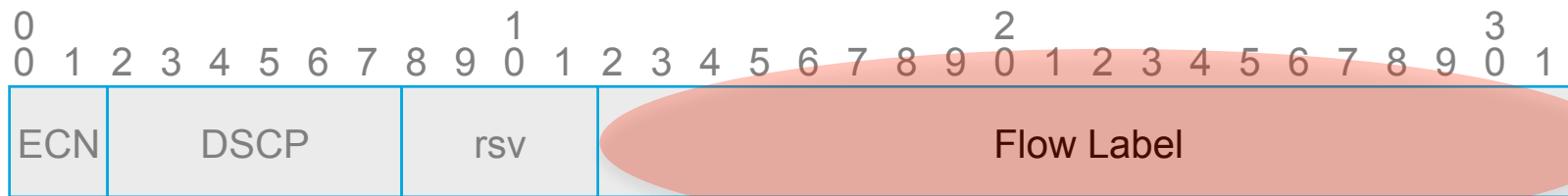


TF	2 bits	Traffic Class and Flow Label
NH	1 bit	Next Header
HLIM	2 bits	Hop Limit
CID	1 bit	Context Identifier Extension
SAC	1 bit	Source Address Context
SAM	2 bits	Source Address Mode
M	1 bit	Multicast Address Compression
DAC	1 bit	Destination Address Context
DAM	2 bits	Destination Address Mode

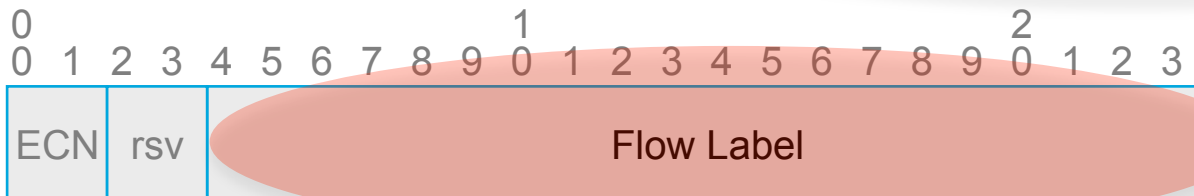
SLoWPAN: Traffic Class & Flow Label



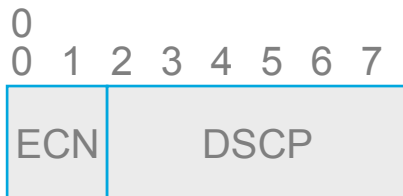
TF = 0



TF = 1



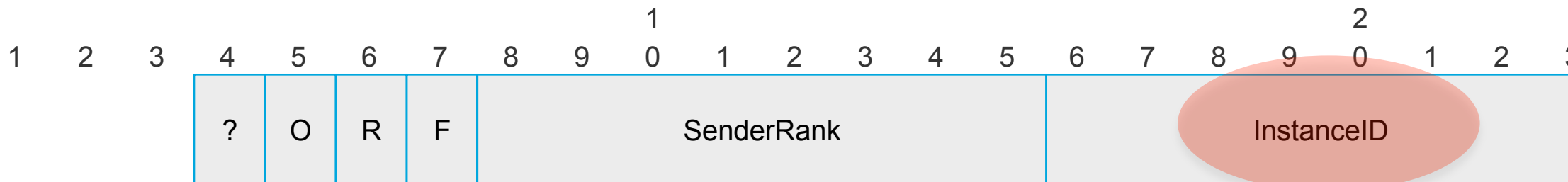
TF = 2



TF = 3

Traffic Class and Flow Label elided.

draft-thubert-roll-flow-label



aces in Flow Label the RPL Packet Information is defined in RFC 6550 [Section 11.2](#)

ave extra HbH header bytes incurred in RFC 6553 AND eventual IPinIP tunneling

“When the router is the source of the original packet and the destination is known to be within the same RPL Instance, the router SHOULD include the RPL Option directly within the original packet. Otherwise, routers MUST use IPv6-in-IPv6 tunneling [\[RFC2473\]](#) and place the RPL Option in the tunnel header.”

discussed with Brian Carpenter on the ROLL ML

<http://www.ietf.org/mail-archive/web/roll/current/msg06967.html>

eed to take to 6MAN to resolve violation of RFC6437

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/OT Network Convergence

Control anything from the network

Learn from Industrial.

Replicate and generalize with open standards

A converged network provides

high availability, flow isolation, security,

Guaranteed bandwidth and **Determinism**

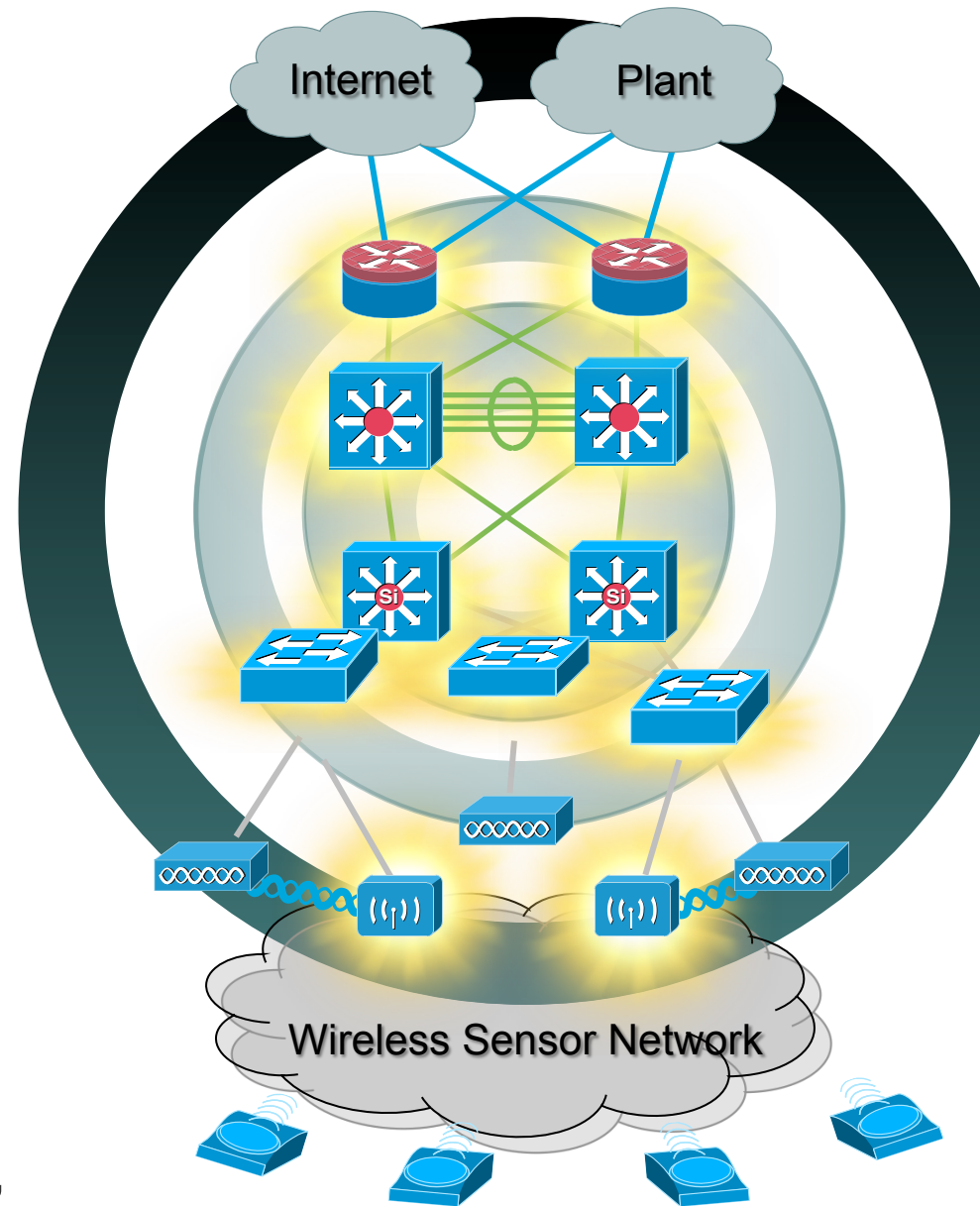
Coexistence involves:

Location and Scheduling

At L2 and L3, for wired and wireless.

New, higher end paradigm

Reaching more devices and data, farther, cheaper,
with better guarantees



Thank you.

